

<http://www.qubit.it>

QUIT
quantum information
theory group

Toward an operational derivation of Quantum Mechanics

Giacomo Mauro D'Ariano

Pavia University

Quantum Physics and Logic, April 9 2009, Oxford Univ. UK

arXiv:0807.4383: in *Philosophy of Quantum Information and Entanglement*, Eds A. Bokulich and G. Jaeger (Cambridge University Press, Cambridge UK, in press)

QM is a probabilistic
theory + *some “principle”*

QM is a probabilistic
theory + *some “principle”*

A principle of the Quantumness,
as the *principle of Relativity*

Quantum
Mechanics

Lorentz
transformations

Quantum
Mechanics

Lorentz
transformations

Principle of
relativity

Existence of a
limiting velocity



Quantum
Mechanics



Principle(s) of
the
Quantumness

Lorentz
transformations



Principle of
relativity

Existence of a
limiting velocity

Quantum
Mechanics



Principle(s) of
the
Quantumness

Lorentz
transformations



Principle of
relativity

Existence of a
limiting velocity



Operational principles indispensable
for local knowability and controllability

Quantum
Mechanics

Principle(s) of
the
Quantumness

Operational principles indispensable
for local knowability and controllability

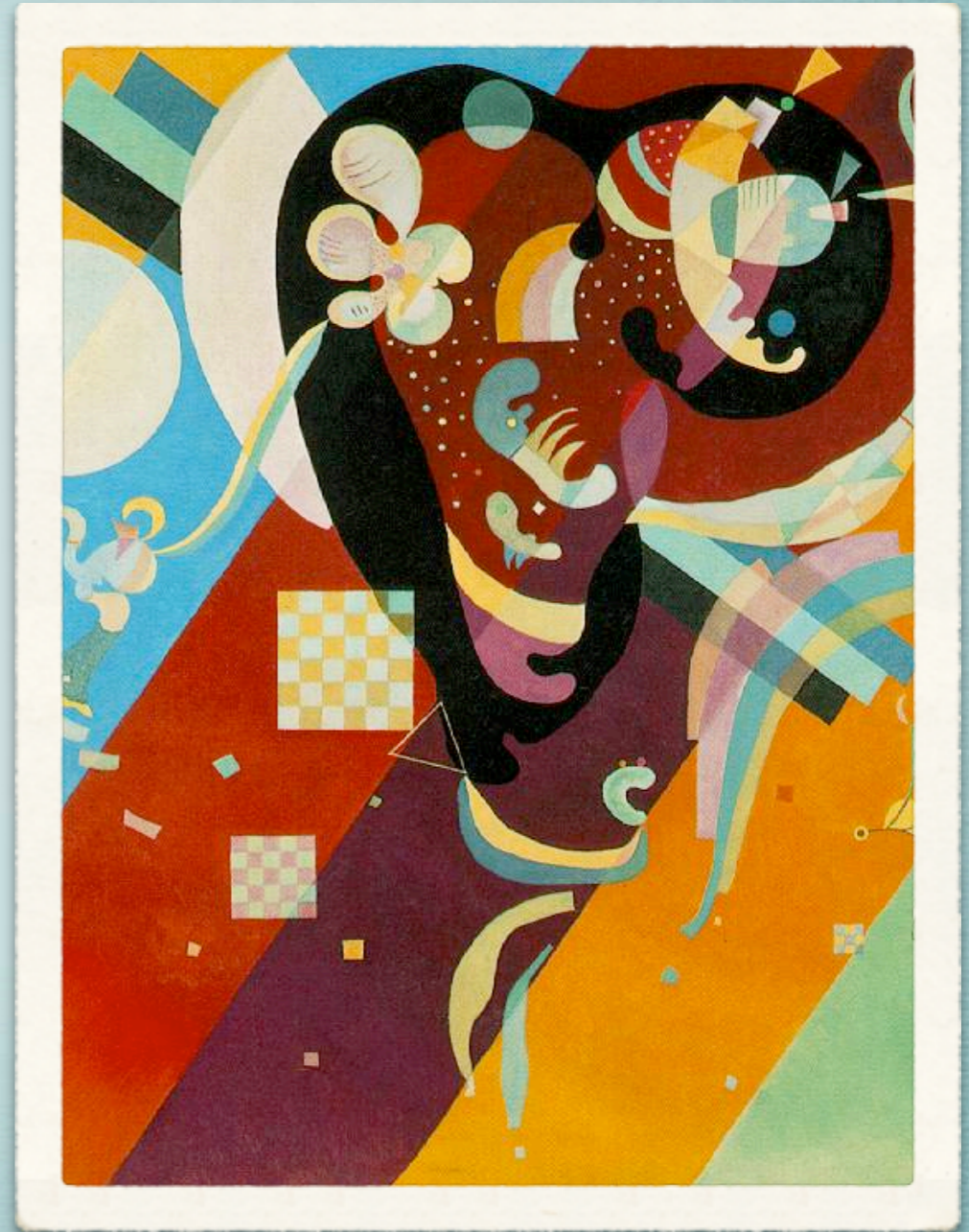
Lorentz
transformations

Principle of
relativity

Existence of a
limiting velocity

In this talk w.l.g.

- * finite dimensions
- * only one kind of system



Operational framework

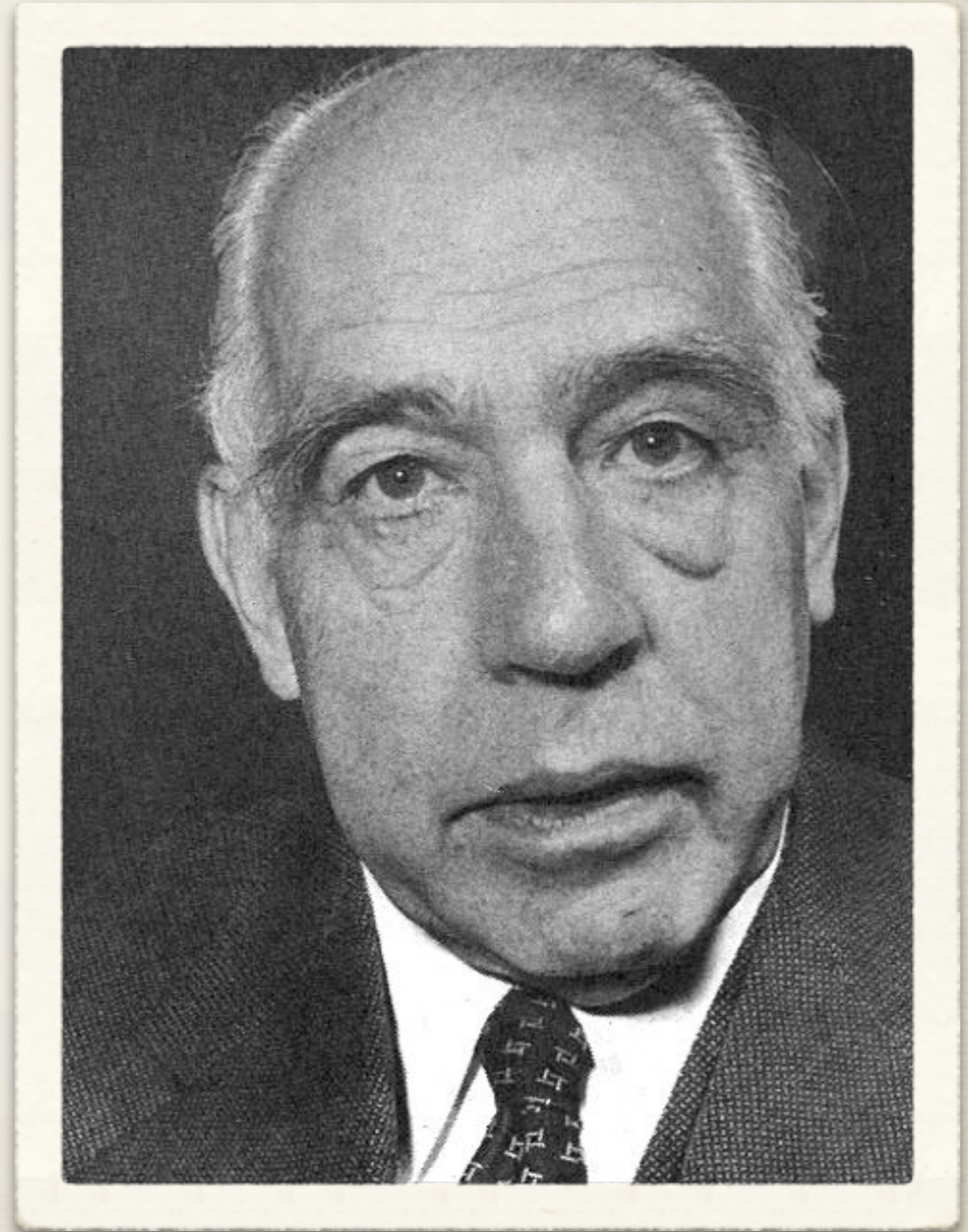
PRIMITIVE NOTIONS

- * probability
- * events
- * ...



General principles

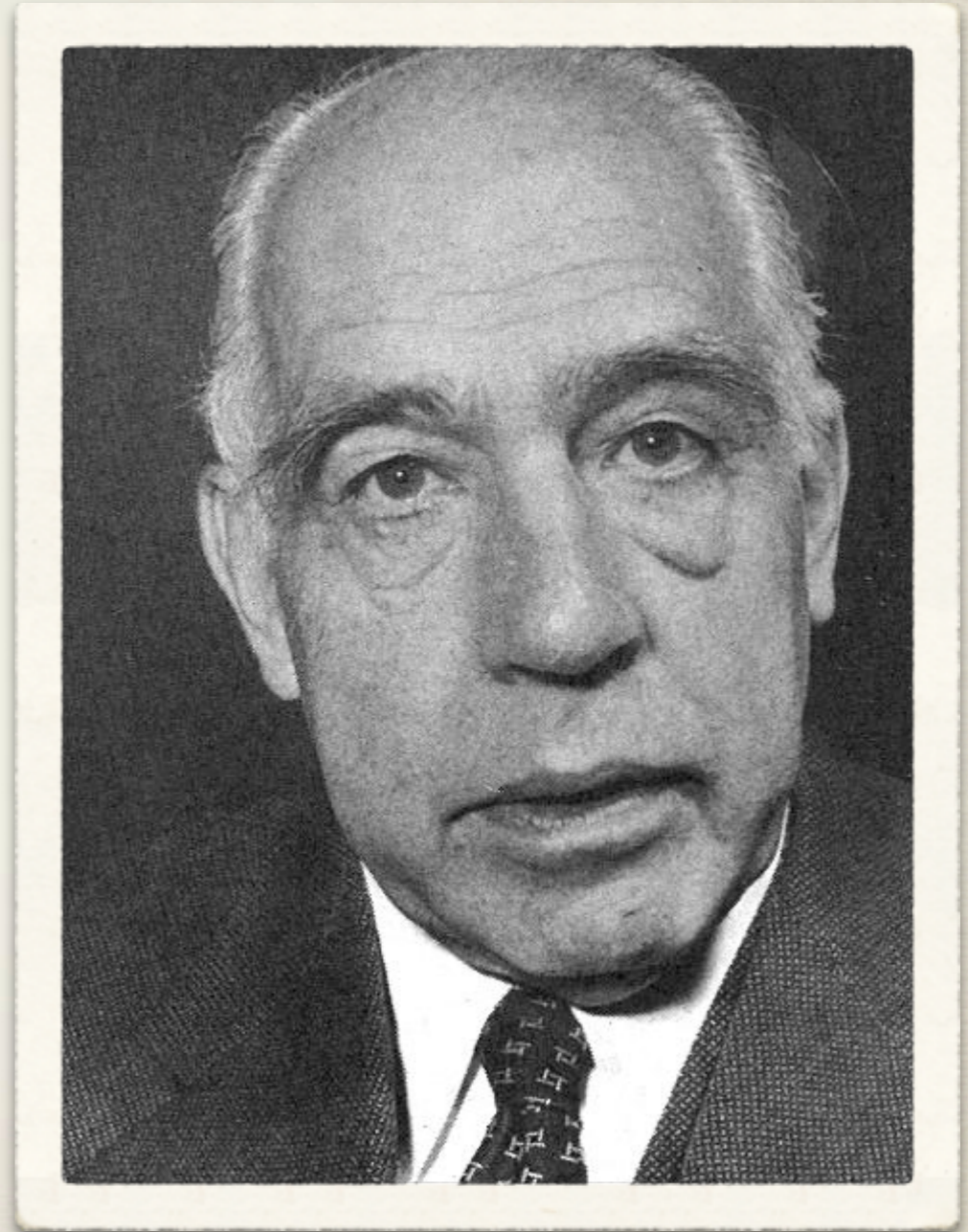
STRONGLY COPENHAGEN



General principles

STRONGLY COPENHAGEN

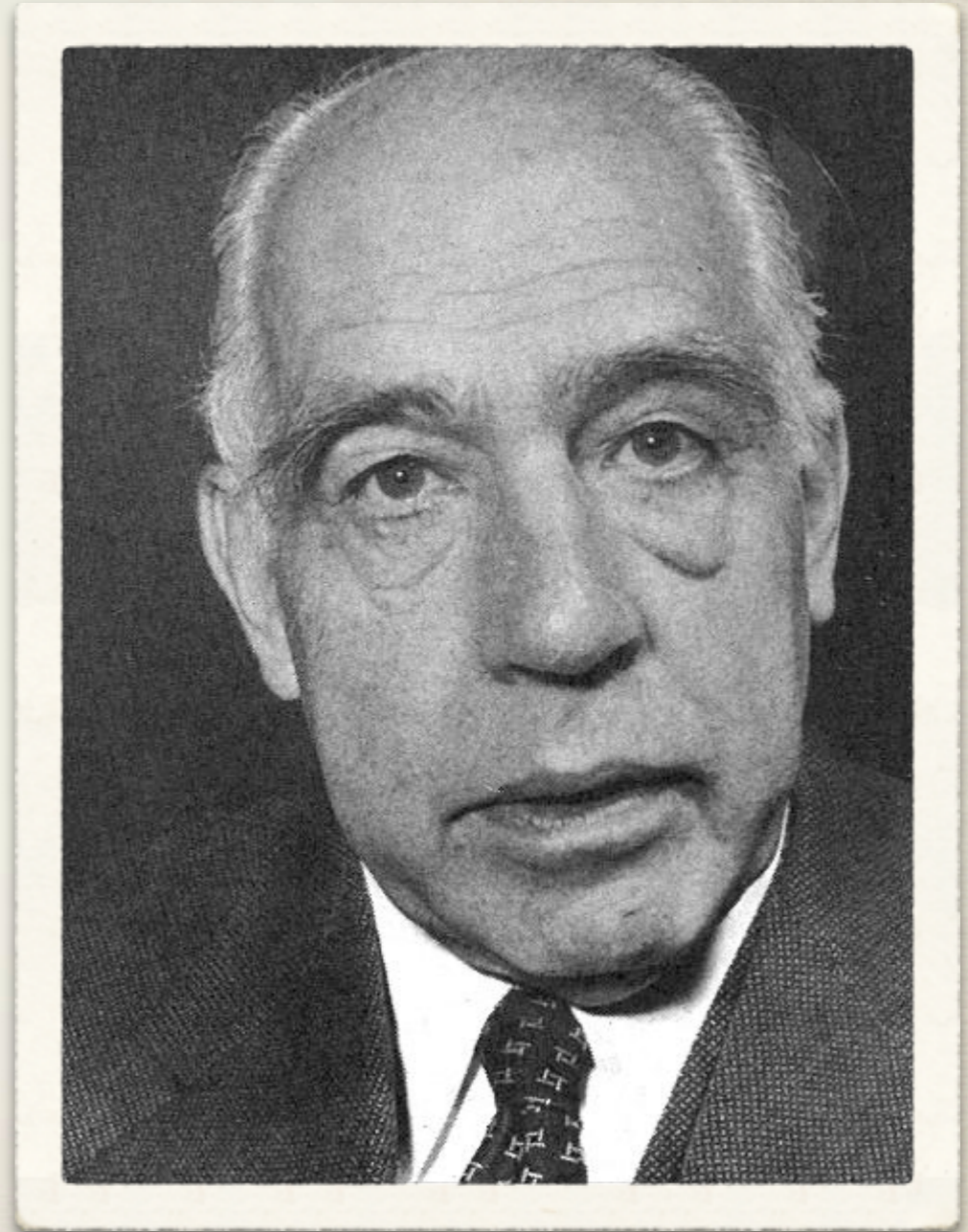
- * everything is defined operationally, including all mathematical objects



General principles

STRONGLY COPENHAGEN

- * everything is defined operationally, including all mathematical objects
- * operational indistinguishability = identification



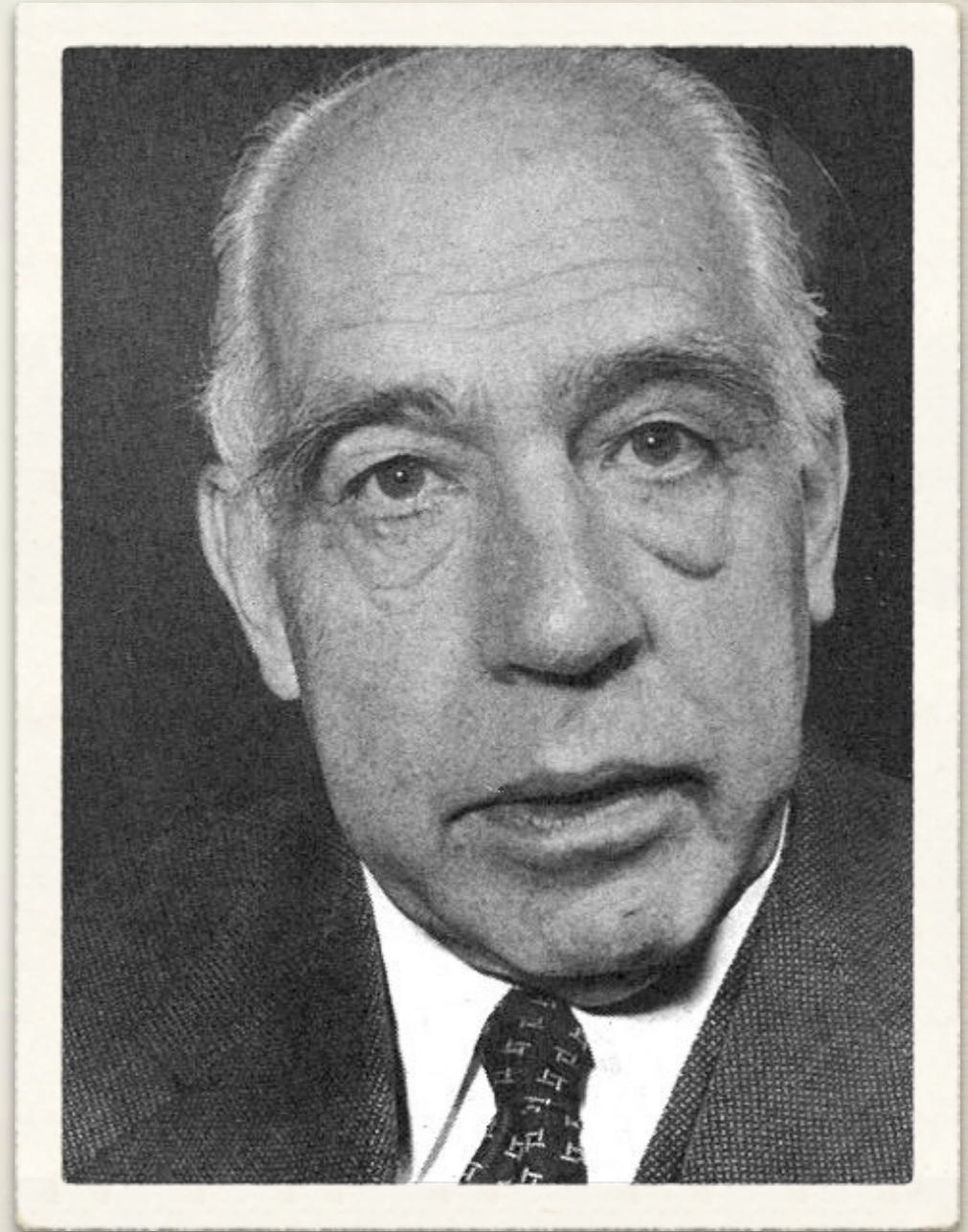
General principles

STRONGLY COPENHAGEN

- * everything is defined operationally, including all mathematical objects
- * operational indistinguishability = identification

Examples:

- * notion of system
- * identification of events
- * sets of states separating effects



General principles

MATHEMATICAL CLOSURE

- * mathematical completion is taken for convenience



General principles

MATHEMATICAL CLOSURE

- * mathematical completion is taken for convenience

Examples:

- * norm closure
- * algebraic closure
- * linear span



General principles

OPERATIONAL CLOSURE

- * every operational option implicit in the formulation is incorporated in the framework



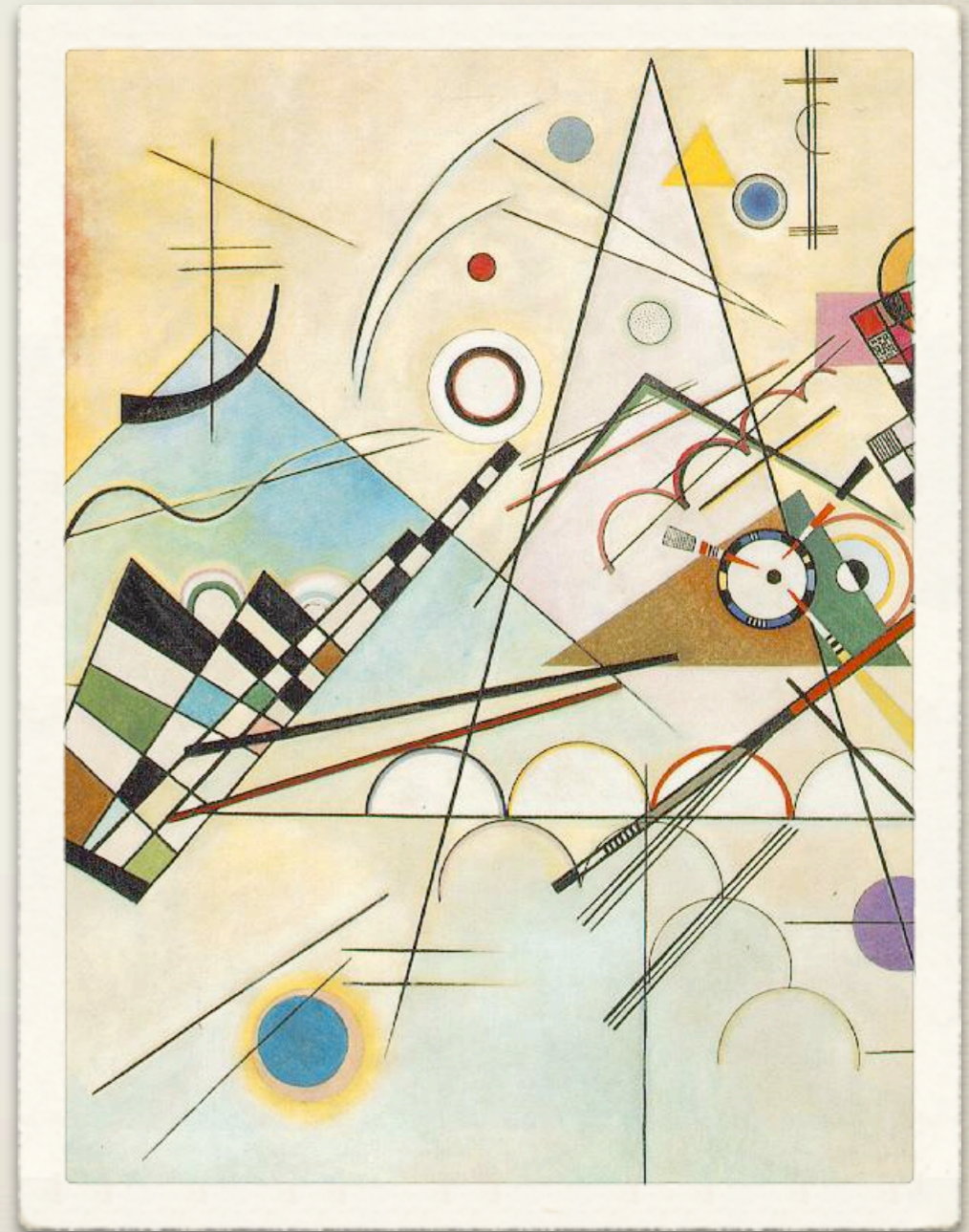
General principles

OPERATIONAL CLOSURE

- * every operational option implicit in the formulation is incorporated in the framework

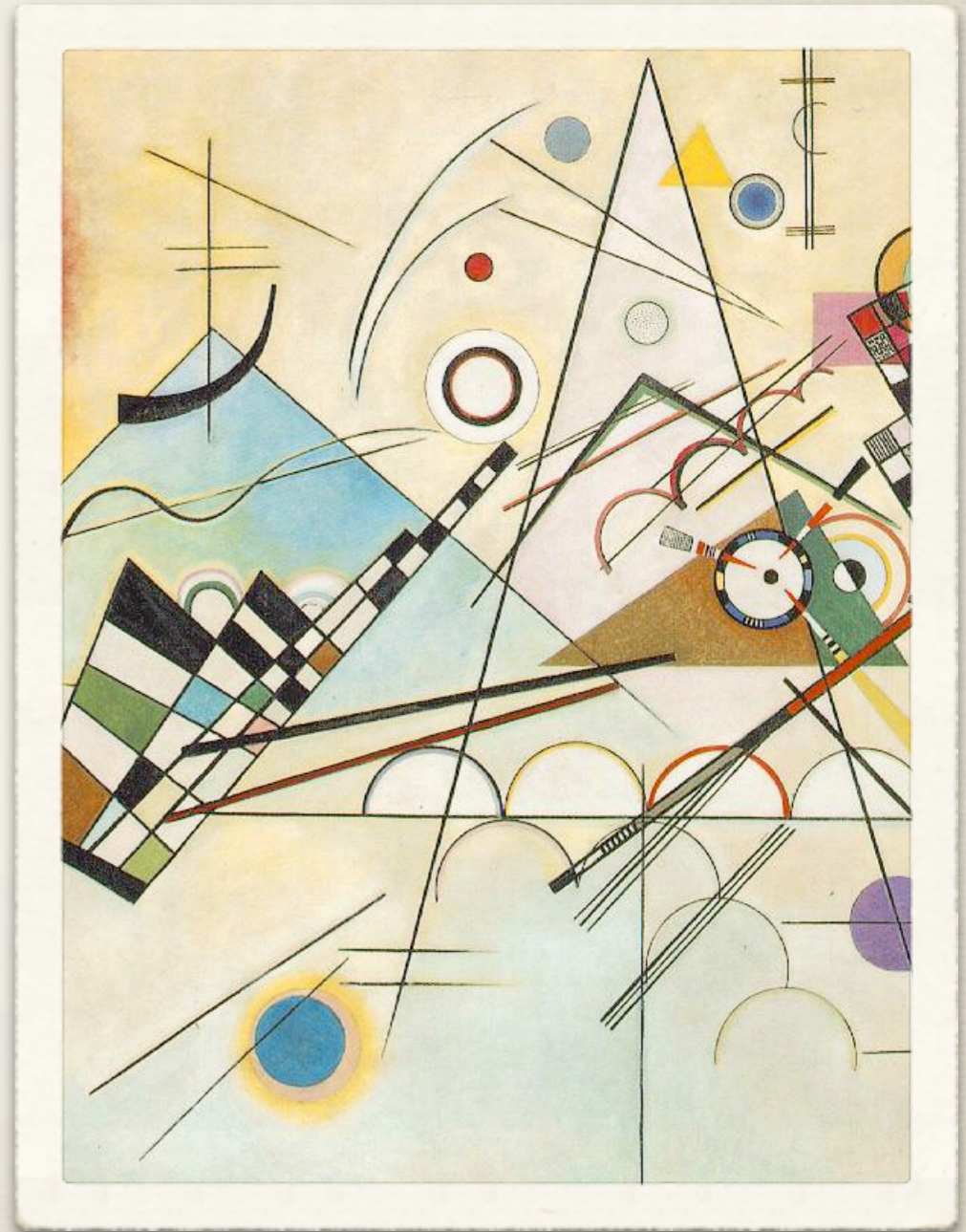
Examples:

- * convex closure
- * closure under coarse-graining



Postulates

- * NSF: **No signaling from the future** (=definition of *cascade*)
- * NS: **No signaling** (=definition of *independent systems*)
- * PFAITH: There exists preparationally faithful states



Postulates under exploration

- * FAITHE: There exists a faithful effect
- * PURIFY: There exists a purification for each state
- * SUPER-PFAITH



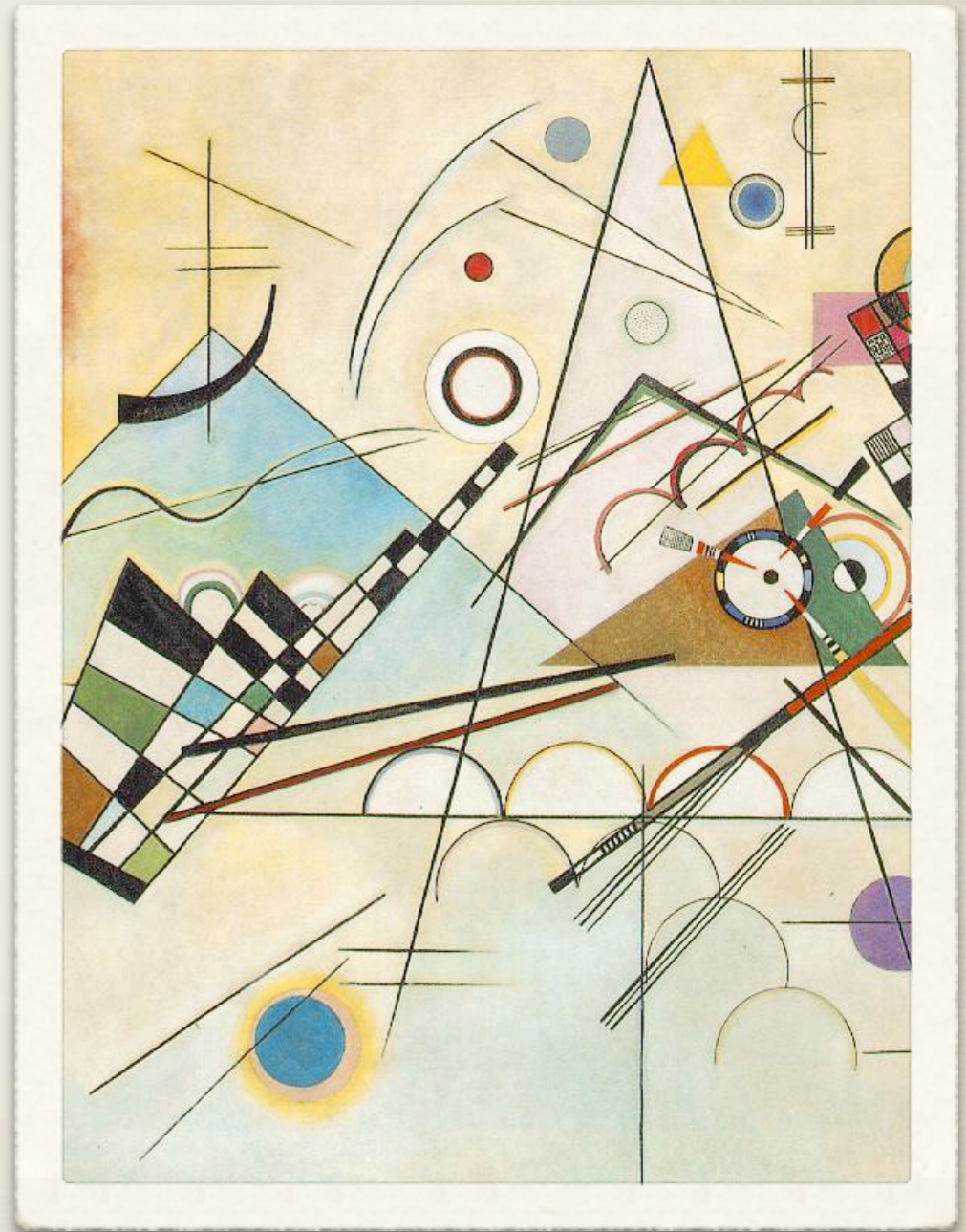
Postulates under exploration

* FAITHE: There exists a faithful effect

* PURIFY: There exists a purification for each state

* SUPER-PFAITH

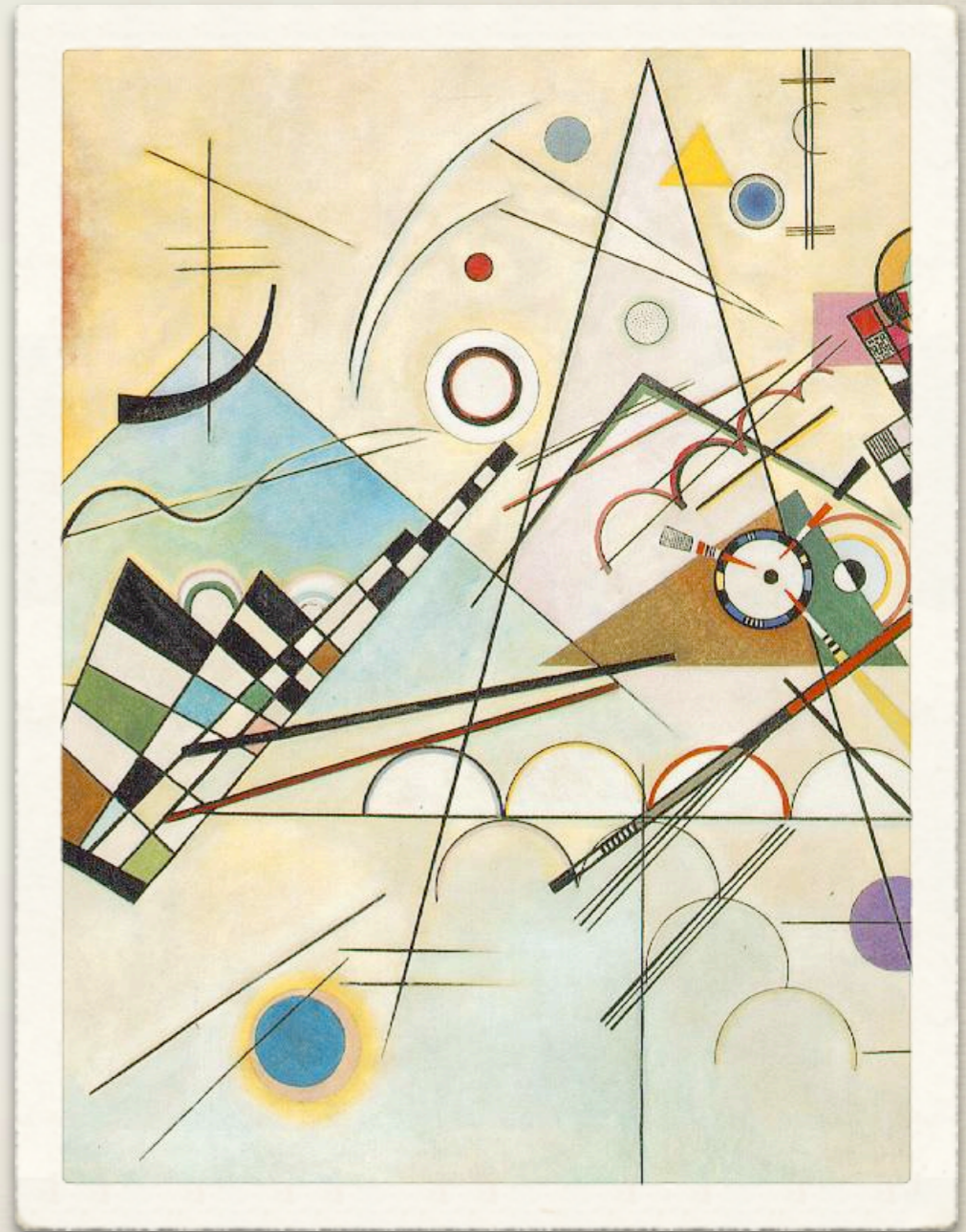
→ Stinespring dilation



Reconstructing QM from probabilities

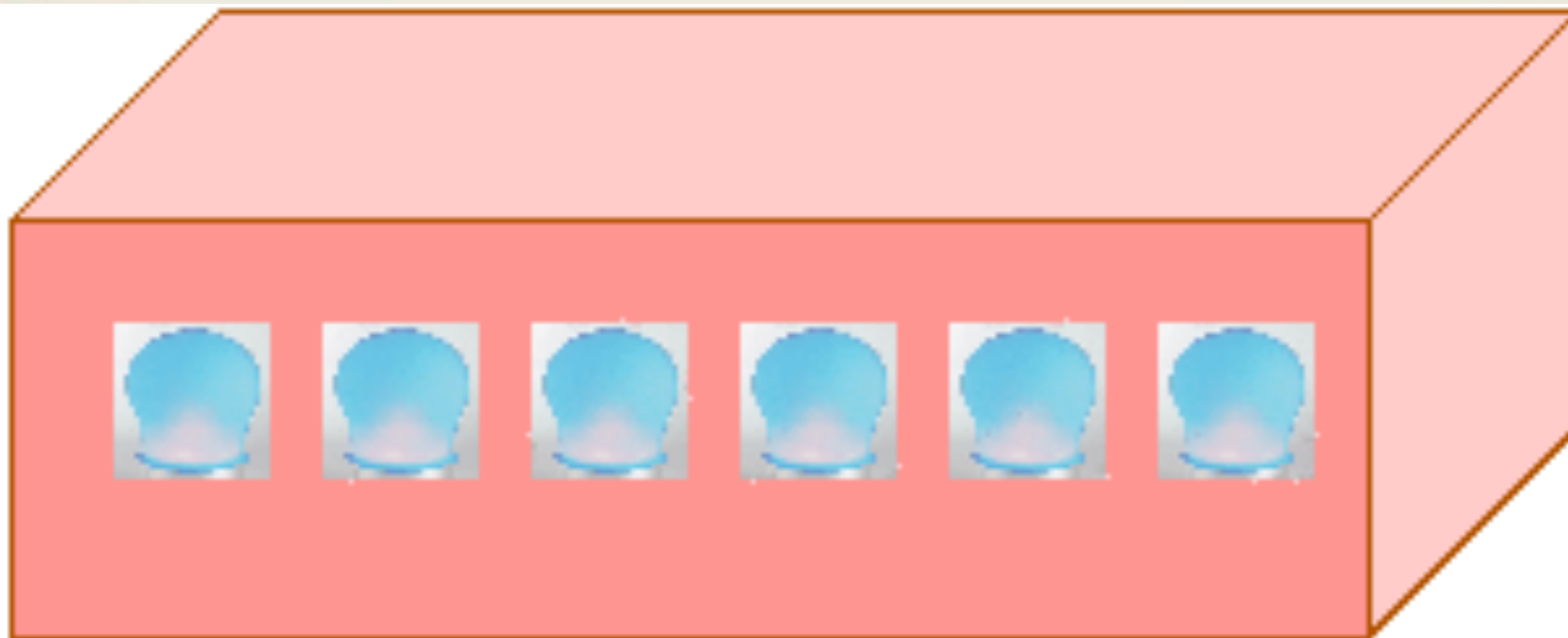
Algebra of effects \leftrightarrow

- * AE: Atomicity of evolution
- * CJ: Choi-Jamiolkowski isomorphism



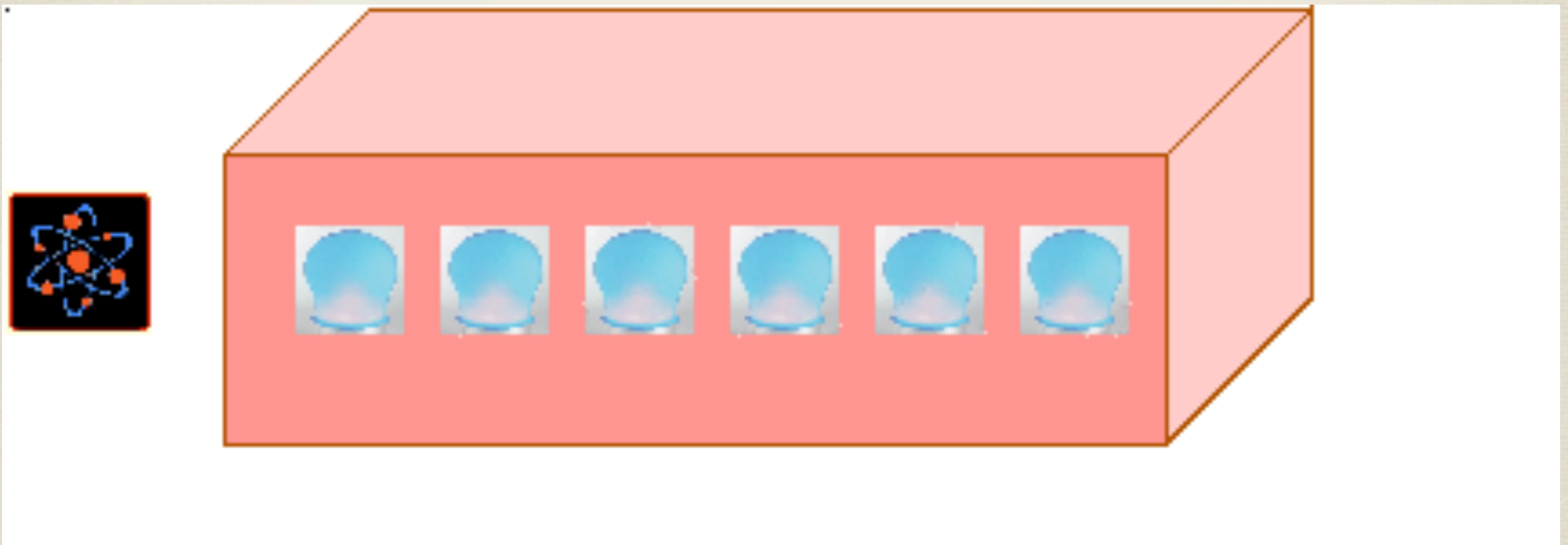
TESTS

• Test/experiment: $\Lambda \equiv \{\mathcal{A}_j\}$ set of possible events \mathcal{A}_j



TESTS

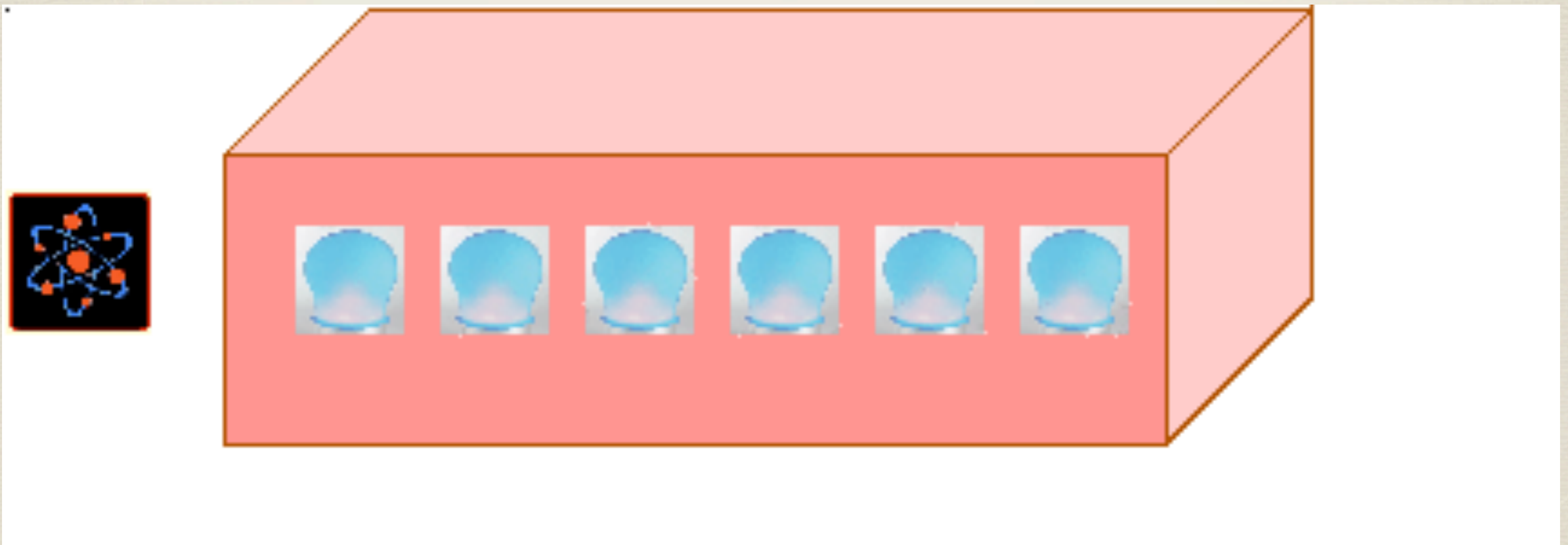
• Test/experiment: $\Lambda \equiv \{\mathcal{A}_j\}$ set of possible events \mathcal{A}_j



(deterministic test/transformation: $\mathbb{D} = \{\mathcal{D}\}$)

TESTS

• Test/experiment: $\Lambda \equiv \{\mathcal{A}_j\}$ set of possible events \mathcal{A}_j



(deterministic test/transformation: $\mathbb{D} = \{\mathcal{D}\}$)

Notice: the same event can occur in different tests

TESTS

Unions of events: $A \cup B$

TESTS

Unions of events: $\mathcal{A} \cup \mathcal{B}$

$$\mathcal{D}_A := \bigcup_{\mathcal{A}_i \in \mathcal{A}} \mathcal{A}_i$$

TESTS

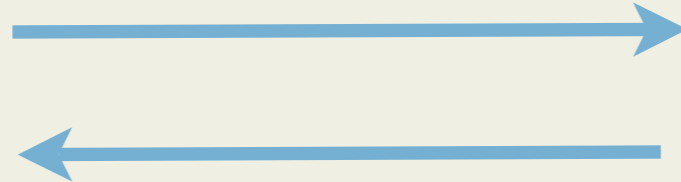
Unions of events: $\mathcal{A} \cup \mathcal{B}$

$$\mathcal{D}_{\mathbb{A}} := \bigcup_{\mathcal{A}_i \in \mathbb{A}} \mathcal{A}_i$$



Coarse-graining

$$\mathbb{A} = \{\mathcal{A}_1, \mathcal{A}_2, \mathcal{A}_3\}$$



$$\mathbb{A}' = \{\mathcal{A}_1, \mathcal{A}_2 \cup \mathcal{A}_3\}$$



Refinement

STATES

State ω : probability rule $\omega(\mathcal{A})$ for any possible event \mathcal{A} in any test

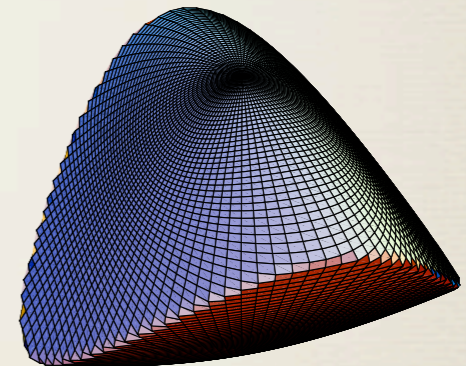
Normalization:
$$\sum_{\mathcal{A}_j \in \mathcal{A}} \omega(\mathcal{A}_j) = 1$$

STATES

State ω : probability rule $\omega(\mathcal{A})$ for any possible event \mathcal{A} in any test

Normalization:
$$\sum_{\mathcal{A}_j \in \mathcal{A}} \omega(\mathcal{A}_j) = 1$$

Convex set of states: \mathcal{S}

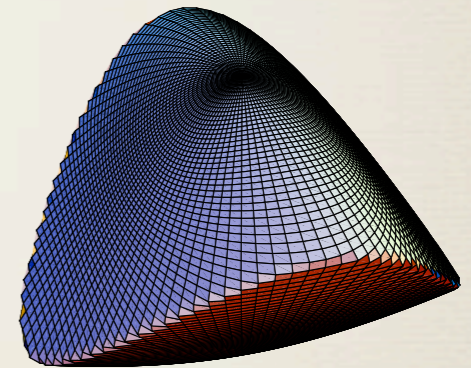


STATES

State ω : probability rule $\omega(\mathcal{A})$ for any possible event \mathcal{A} in any test

Normalization:
$$\sum_{\mathcal{A}_j \in \mathcal{A}} \omega(\mathcal{A}_j) = 1$$

Convex set of states: \mathcal{S}

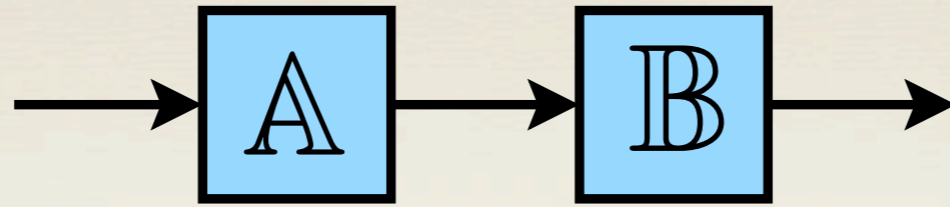


States will also be regarded as tests themselves:
“preparation-tests”.

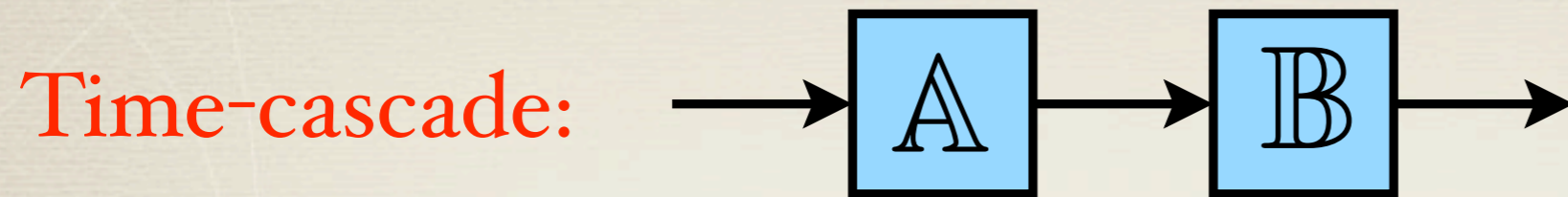
CASCADES OF TESTS

CASCADES OF TESTS

Time-cascade:



CASCADES OF TESTS



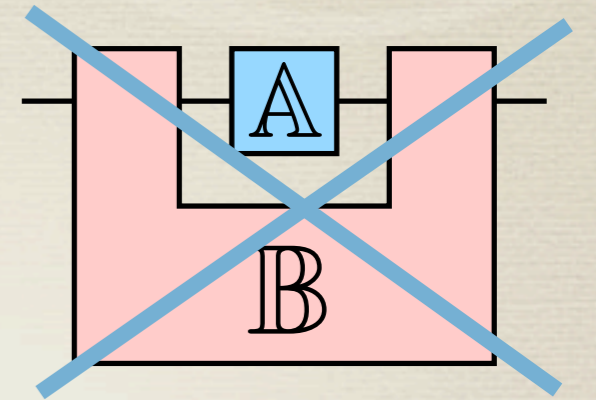
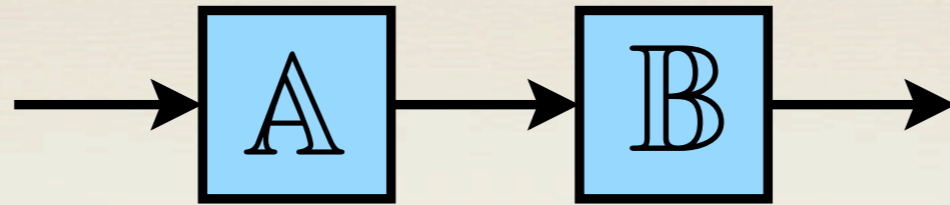
$$\mathbb{B} \circ \mathbb{A} = \{\mathcal{B}_j \circ \mathcal{A}_i\} \text{ cascade of tests } \mathbb{A} = \{\mathcal{A}_i\}, \mathbb{B} = \{\mathcal{B}_j\},$$

collection of joined events with the following rule for marginals:

$$\mathbf{NSF} \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) =: f(\mathbb{B}, \mathcal{A}) \equiv \omega(\mathcal{A}), \quad \forall \mathbb{B}, \mathcal{A}, \omega$$

CASCADES OF TESTS

Time-cascade:



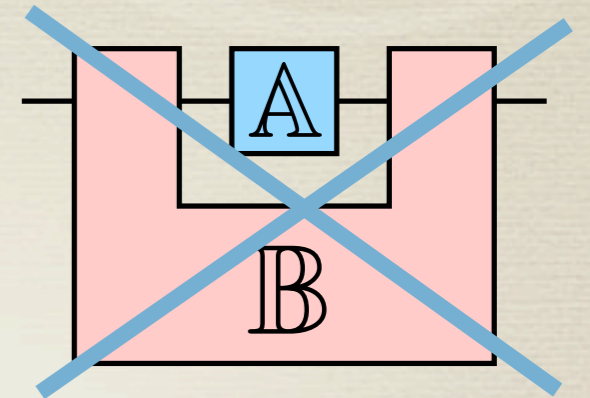
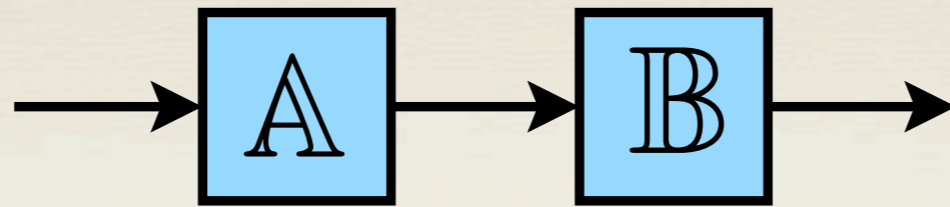
$\mathbb{B} \circ \mathbb{A} = \{\mathcal{B}_j \circ \mathcal{A}_i\}$ cascade of tests $\mathbb{A} = \{\mathcal{A}_i\}$, $\mathbb{B} = \{\mathcal{B}_j\}$,

collection of joined events with the following rule for marginals:

$$\mathbf{NSF} \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) =: f(\mathbb{B}, \mathcal{A}) \equiv \omega(\mathcal{A}), \quad \forall \mathbb{B}, \mathcal{A}, \omega$$

CASCADES OF TESTS

Time-cascade:



$$\mathbb{B} \circ \mathbb{A} = \{\mathcal{B}_j \circ \mathcal{A}_i\} \text{ cascade of tests } \mathbb{A} = \{\mathcal{A}_i\}, \mathbb{B} = \{\mathcal{B}_j\},$$

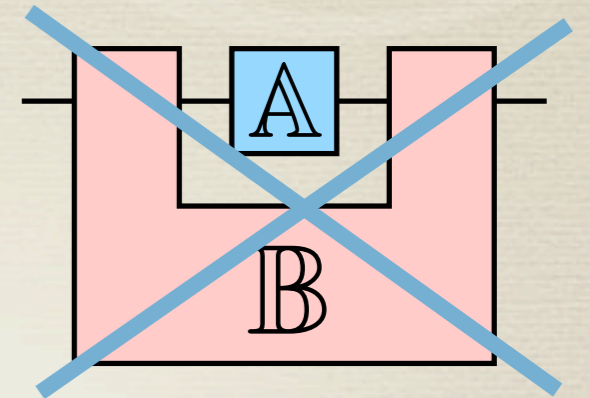
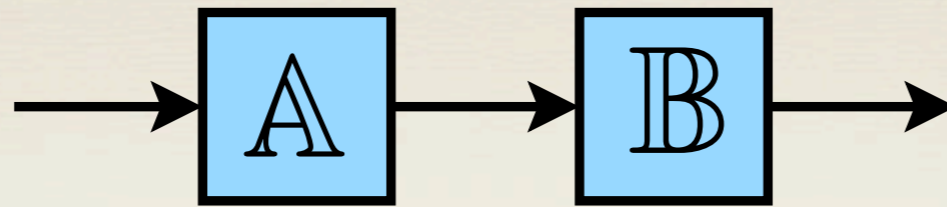
collection of joined events with the following rule for marginals:

$$\mathbf{NSF} \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) =: f(\mathbb{B}, \mathcal{A}) \equiv \omega(\mathcal{A}), \quad \forall \mathbb{B}, \mathcal{A}, \omega$$

➔ composition of events: $\mathcal{B} \circ \mathcal{A}$

CASCADES OF TESTS

Time-cascade:



$$\mathbb{B} \circ \mathbb{A} = \{\mathcal{B}_j \circ \mathcal{A}_i\} \text{ cascade of tests } \mathbb{A} = \{\mathcal{A}_i\}, \mathbb{B} = \{\mathcal{B}_j\},$$

collection of joined events with the following rule for marginals:

$$\mathbf{NSF} \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) =: f(\mathbb{B}, \mathcal{A}) \equiv \omega(\mathcal{A}), \quad \forall \mathbb{B}, \mathcal{A}, \omega$$



composition of events: $\mathcal{B} \circ \mathcal{A}$

Convex monoid
of events:

\mathfrak{T}

Events \equiv transformations

$$\text{NSF} \quad \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) = \omega(\mathcal{A}), \quad \forall \mathbb{B}, \forall \mathcal{A}, \forall \omega$$

Events \equiv transformations

$$\text{NSF} \quad \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) = \omega(\mathcal{A}), \quad \forall \mathbb{B}, \forall \mathcal{A}, \forall \omega$$

\Rightarrow conditional probability: $p(\mathcal{B}|\mathcal{A}) = \omega(\mathcal{B} \circ \mathcal{A})/\omega(\mathcal{A})$

\Rightarrow conditional state: $\omega_{\mathcal{A}} := \omega(\cdot \circ \mathcal{A})/\omega(\mathcal{A})$

variable

Events \equiv transformations

$$\text{NSF} \quad \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) = \omega(\mathcal{A}), \quad \forall \mathbb{B}, \forall \mathcal{A}, \forall \omega$$

\Rightarrow conditional probability: $p(\mathcal{B}|\mathcal{A}) = \omega(\mathcal{B} \circ \mathcal{A}) / \omega(\mathcal{A})$

\Rightarrow conditional state: $\omega_{\mathcal{A}} := \omega(\cdot \circ \mathcal{A}) / \omega(\mathcal{A})$

\Rightarrow evolution \equiv state conditioning: $\mathcal{A}\omega := \omega(\cdot \circ \mathcal{A})$

\Rightarrow events \equiv transformations

Events \equiv transformations

$$\text{NSF} \quad \sum_{\mathcal{B}_j \in \mathbb{B}} \omega(\mathcal{B}_j \circ \mathcal{A}) = \omega(\mathcal{A}), \quad \forall \mathbb{B}, \forall \mathcal{A}, \forall \omega$$

\Rightarrow conditional probability: $p(\mathcal{B}|\mathcal{A}) = \omega(\mathcal{B} \circ \mathcal{A}) / \omega(\mathcal{A})$

\Rightarrow conditional state: $\omega_{\mathcal{A}} := \omega(\cdot \circ \mathcal{A}) / \omega(\mathcal{A})$

\Rightarrow evolution \equiv state conditioning: $\mathcal{A}\omega := \omega(\cdot \circ \mathcal{A})$

\Rightarrow events \equiv transformations

\Rightarrow linearity of evolution

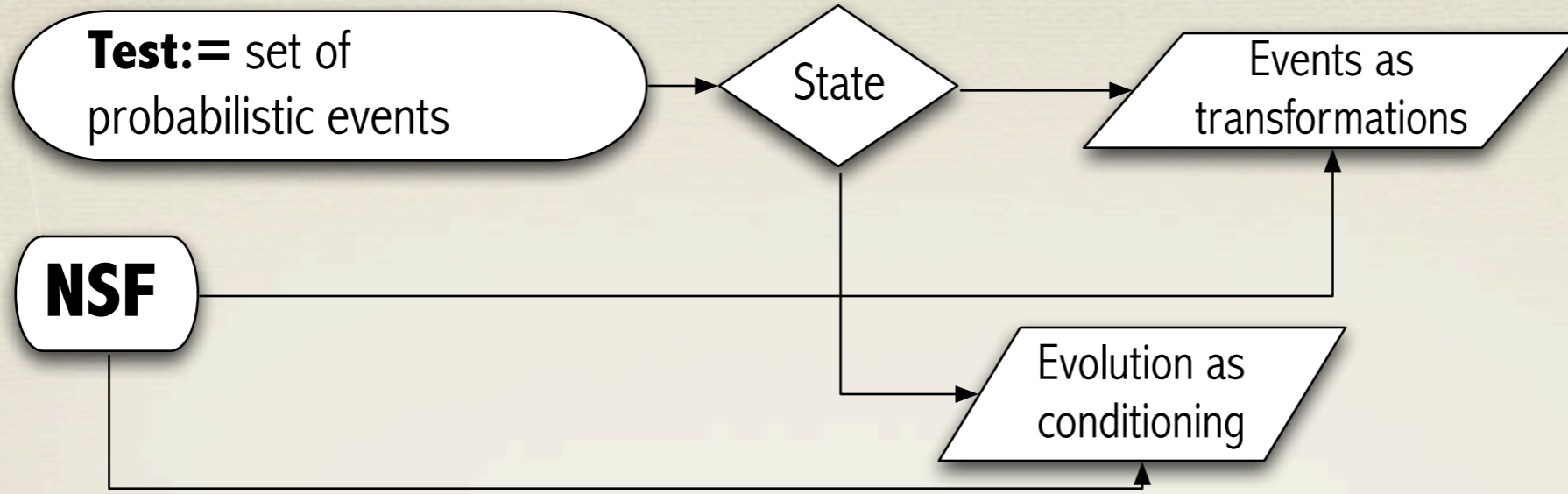
Logical flow-diagram

Test:= set of
probabilistic events

State

```
graph LR; A([Test:= set of probabilistic events]) --> B{State}
```

Logical flow-diagram

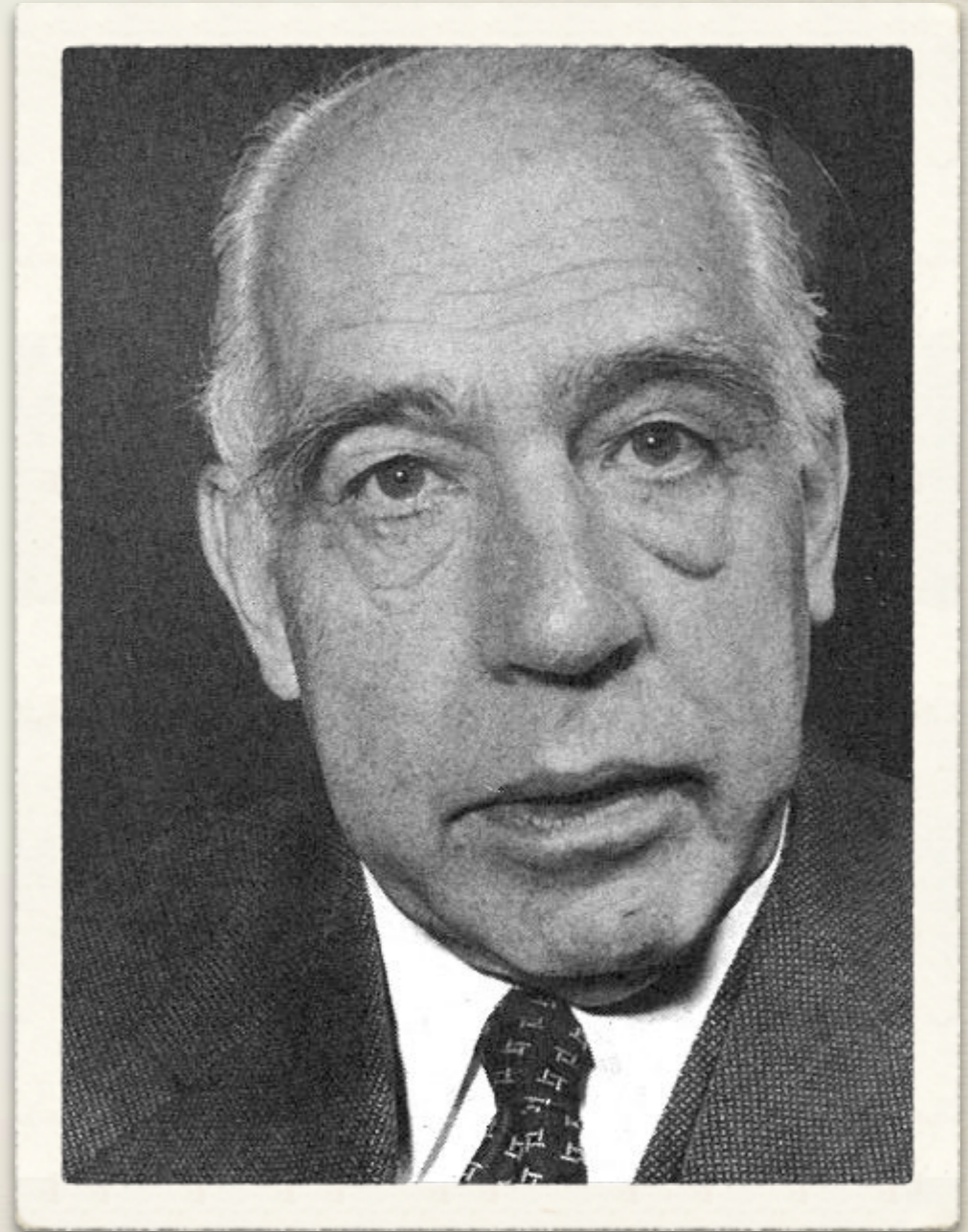
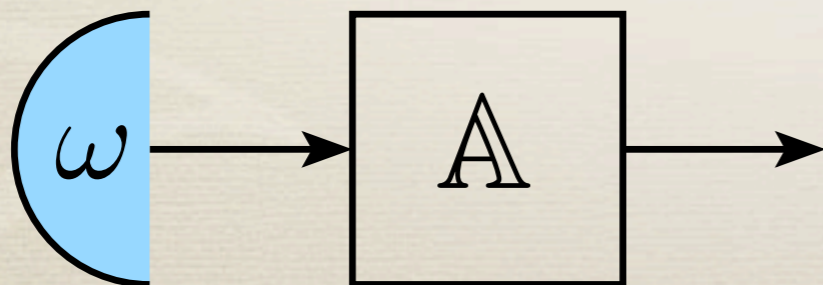


SYSTEM

$$S = \{\omega_1, \omega_2, \dots, A, B, C, \dots\}$$

collection of tests closed under

- * coarse-graining
- * conditioning
- * cascading
- * convex combination



Copenhagen

SYSTEM

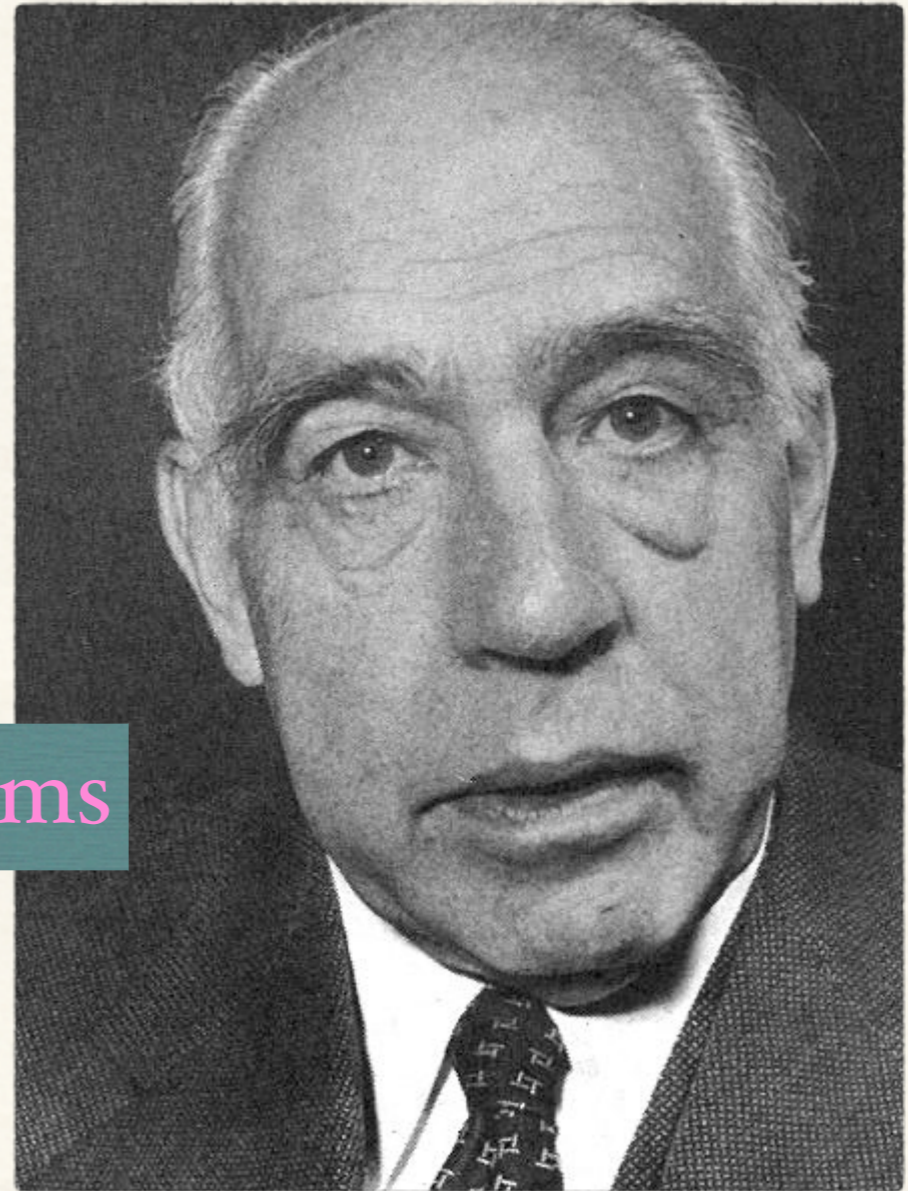
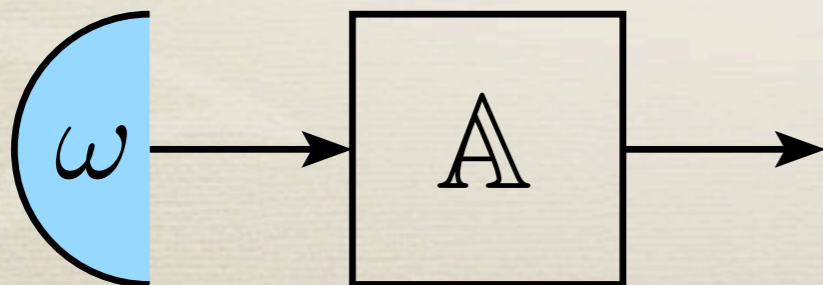
$$S = \{\omega_1, \omega_2, \dots, A, B, C, \dots\}$$

Copenhagen

collection of tests closed under

- * coarse-graining
- * conditioning
- * cascading
- * convex combination

one kind of systems



different systems: tests as letters in a “language”

Two equivalence classes for transformations

Two equivalence classes for transformations

Two transformations \mathcal{A} and \mathcal{B} are
conditioning equivalent if

$$\omega_{\mathcal{A}} = \omega_{\mathcal{B}} \quad \forall \omega \in \mathcal{G}$$

Two equivalence classes for transformations

Two transformations \mathcal{A} and \mathcal{B} are **conditioning equivalent** if

$$\omega_{\mathcal{A}} = \omega_{\mathcal{B}} \quad \forall \omega \in \mathcal{G}$$

Two transformations \mathcal{A} and \mathcal{B} are **probabilistically equivalent** if

$$\omega(\mathcal{A}) = \omega(\mathcal{B}) \quad \forall \omega \in \mathcal{G}$$

Two equivalence classes for transformations

Two transformations \mathcal{A} and \mathcal{B} are **conditioning equivalent** if

$$\omega_{\mathcal{A}} = \omega_{\mathcal{B}} \quad \forall \omega \in \mathcal{G}$$

Conditioning-equivalence class

Two transformations \mathcal{A} and \mathcal{B} are **probabilistically equivalent** if

$$\omega(\mathcal{A}) = \omega(\mathcal{B}) \quad \forall \omega \in \mathcal{G}$$

Probabilistic-equivalence class

Two equivalence classes for transformations

A transformation is completely specified by the two classes:

$$\mathcal{A}\omega = \omega(\mathcal{A})\omega_{\mathcal{A}}$$

↑
probabilistic

↑
conditioning

Two equivalence classes for transformations

A transformation is completely specified by the two classes:

$$\mathcal{A}\omega = \omega(\mathcal{A})\omega_{\mathcal{A}}$$

↑
probabilistic

↑
conditioning

variable
↓

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$

EFFECTS

Effect $[\mathcal{A}]_{\text{eff}}$: equivalence class of transformations occurring with the same probability as \mathcal{A} for all states.

$$\forall \omega \in \mathcal{S} : \quad \omega(\mathcal{A}) \equiv \omega([\mathcal{A}]_{\text{eff}})$$

EFFECTS

Effect $[\mathcal{A}]_{\text{eff}}$: equivalence class of transformations occurring with the same probability as \mathcal{A} for all states.

$$\forall \omega \in \mathcal{S} : \quad \omega(\mathcal{A}) \equiv \omega([\mathcal{A}]_{\text{eff}})$$

a effect $\rightarrow \mathcal{A} \in a$ means $\omega(\mathcal{A}) \equiv \omega(a)$

EFFECTS

Effect $[\mathcal{A}]_{\text{eff}}$: equivalence class of transformations occurring with the same probability as \mathcal{A} for all states.

$$\forall \omega \in \mathcal{S} : \quad \omega(\mathcal{A}) \equiv \omega([\mathcal{A}]_{\text{eff}})$$

a effect $\rightarrow \mathcal{A} \in a$ means $\omega(\mathcal{A}) \equiv \omega(a)$

\mathcal{E} := convex set of effects

EFFECTS

Effect $[\mathcal{A}]_{\text{eff}}$: equivalence class of transformations occurring with the same probability as \mathcal{A} for all states.

$$\forall \omega \in \mathcal{S} : \quad \omega(\mathcal{A}) \equiv \omega([\mathcal{A}]_{\text{eff}})$$

a effect $\rightarrow \mathcal{A} \in a$ means $\omega(\mathcal{A}) \equiv \omega(a)$

$\mathcal{E} :=$ convex set of effects

Duality: effects \mathcal{E} positive linear functionals over states (bounded by 1) $a \in \mathcal{E}, \omega \in \mathcal{S}, \quad \omega(a) \equiv a(\omega)$

EFFECTS

Effect $[\mathcal{A}]_{\text{eff}}$: equivalence class of transformations occurring with the same probability as \mathcal{A} for all states.

$$\forall \omega \in \mathcal{S} : \quad \omega(\mathcal{A}) \equiv \omega([\mathcal{A}]_{\text{eff}})$$

a effect $\rightarrow \mathcal{A} \in a$ means $\omega(\mathcal{A}) \equiv \omega(a)$

\mathcal{E} := convex set of effects

Duality: effects \mathcal{E} positive linear functionals over states (bounded by 1) $a \in \mathcal{E}, \omega \in \mathcal{S}, \quad \omega(a) \equiv a(\omega)$

e **deterministic effect** i.e. $\omega(e) = 1 \quad \forall \omega \in \mathcal{S}$

EFFECTS

EFFECTS

State-conditioning \Rightarrow Transformations act linearly over effects:

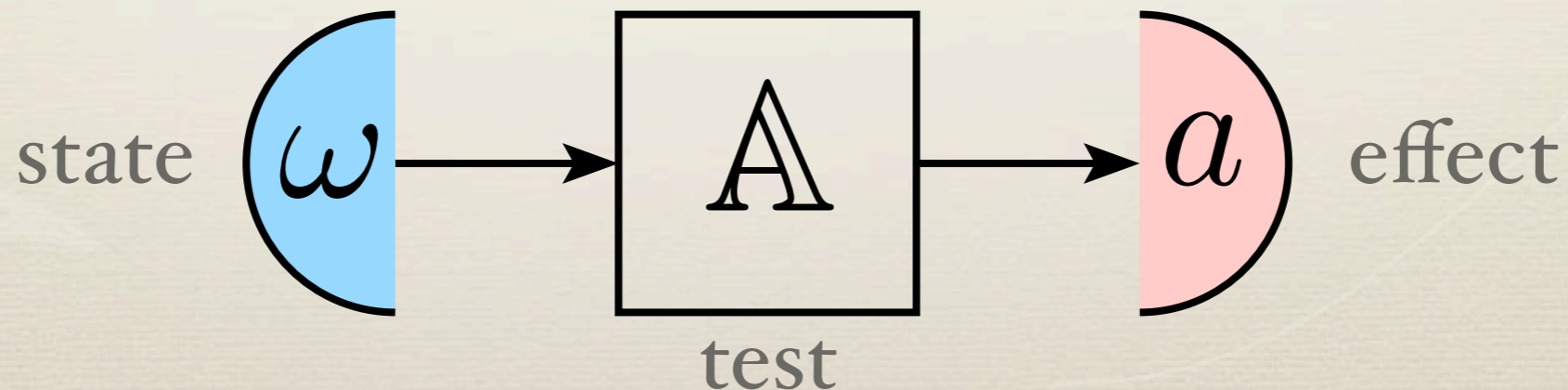
$$[\mathcal{B}]_{\text{eff}} \circ \mathcal{A} = [\mathcal{B} \circ \mathcal{A}]_{\text{eff}} \quad (\text{Heisenberg picture})$$

EFFECTS

State-conditioning \Rightarrow Transformations act linearly over effects:

$$[\mathcal{B}]_{\text{eff}} \circ \mathcal{A} = [\mathcal{B} \circ \mathcal{A}]_{\text{eff}} \quad (\text{Heisenberg picture})$$

Effects will also be regarded as tests themselves: “effect-tests”

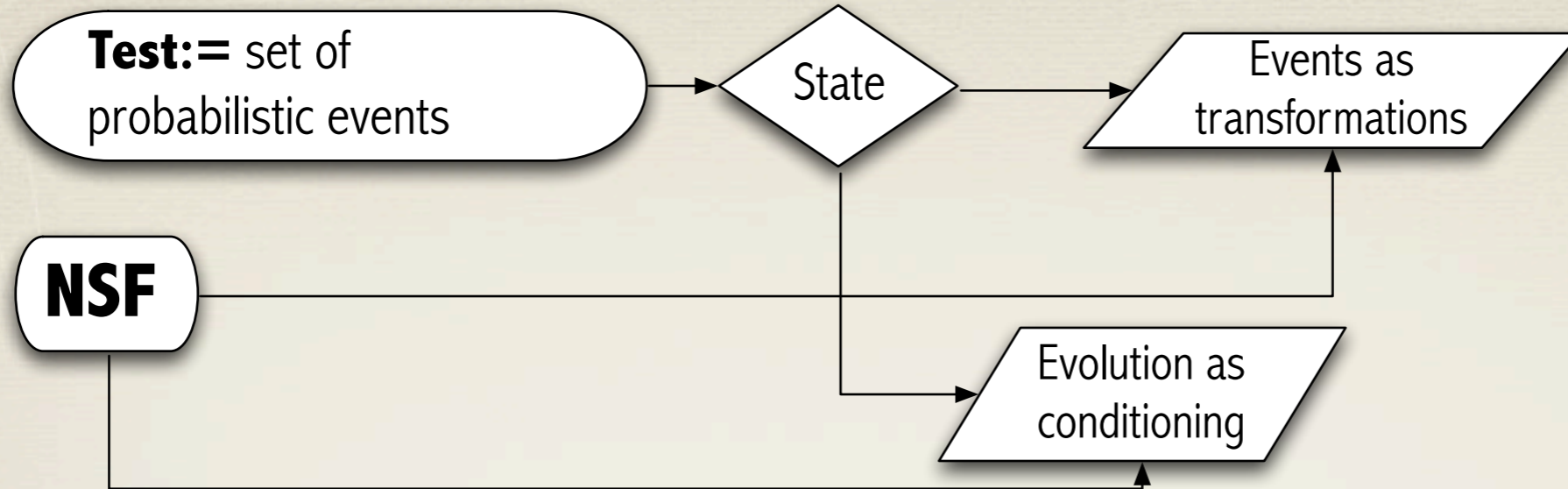


OBSERVABLE

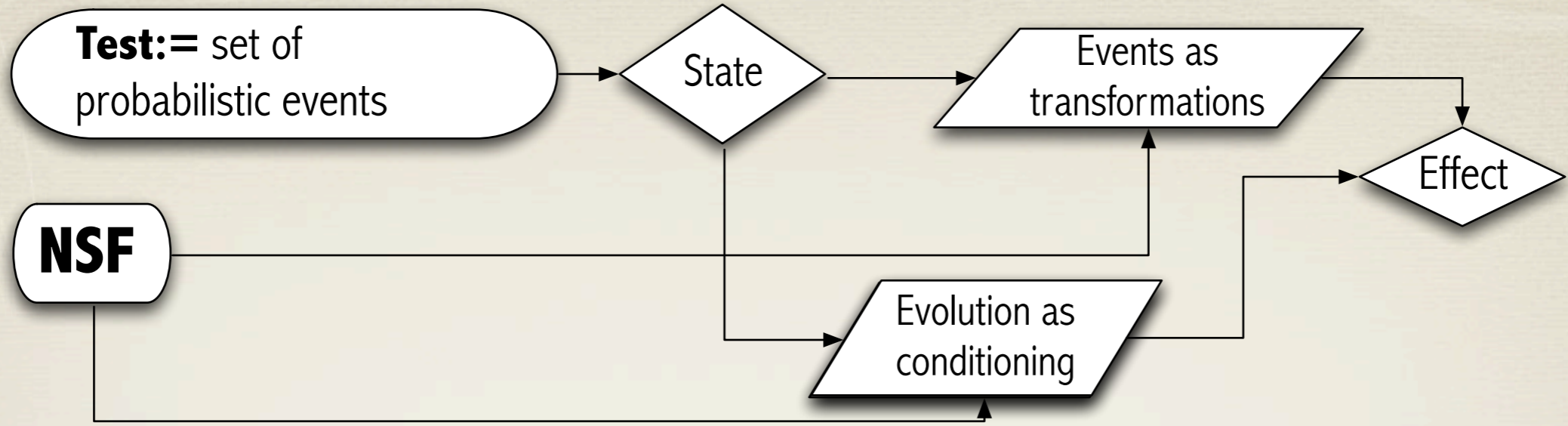
Observable $\mathbb{L} = \{l_i\}$: complete set of effects of a test

Normalization:
$$\sum_{i \in \mathbb{L}} l_i = e$$

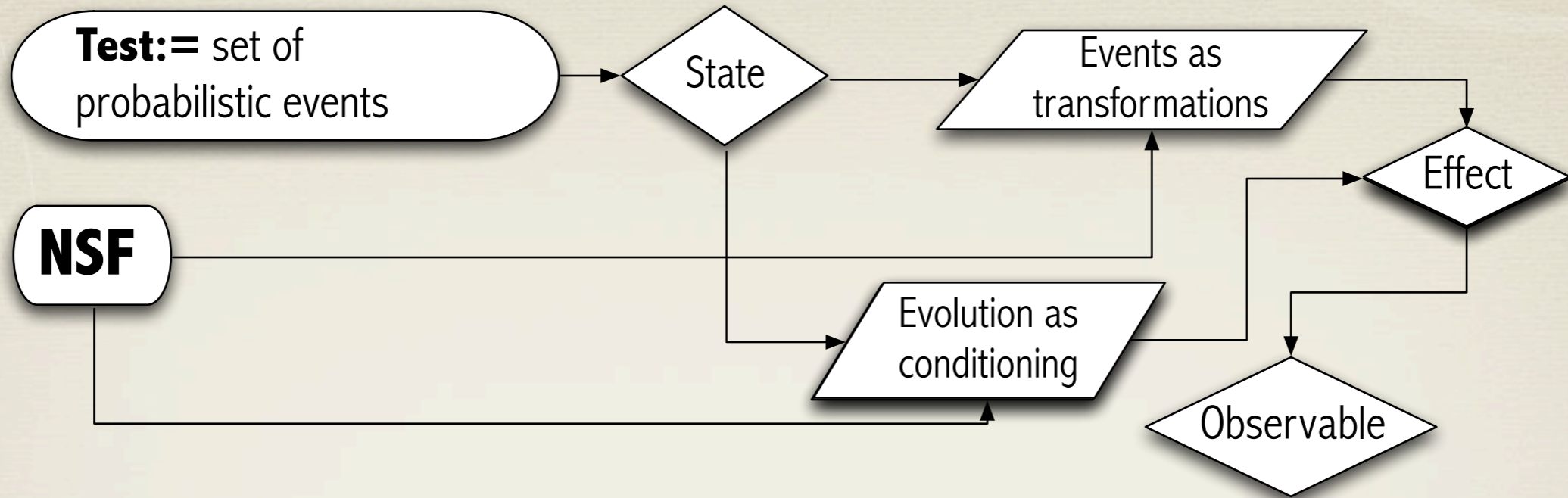
Logical flow-diagram



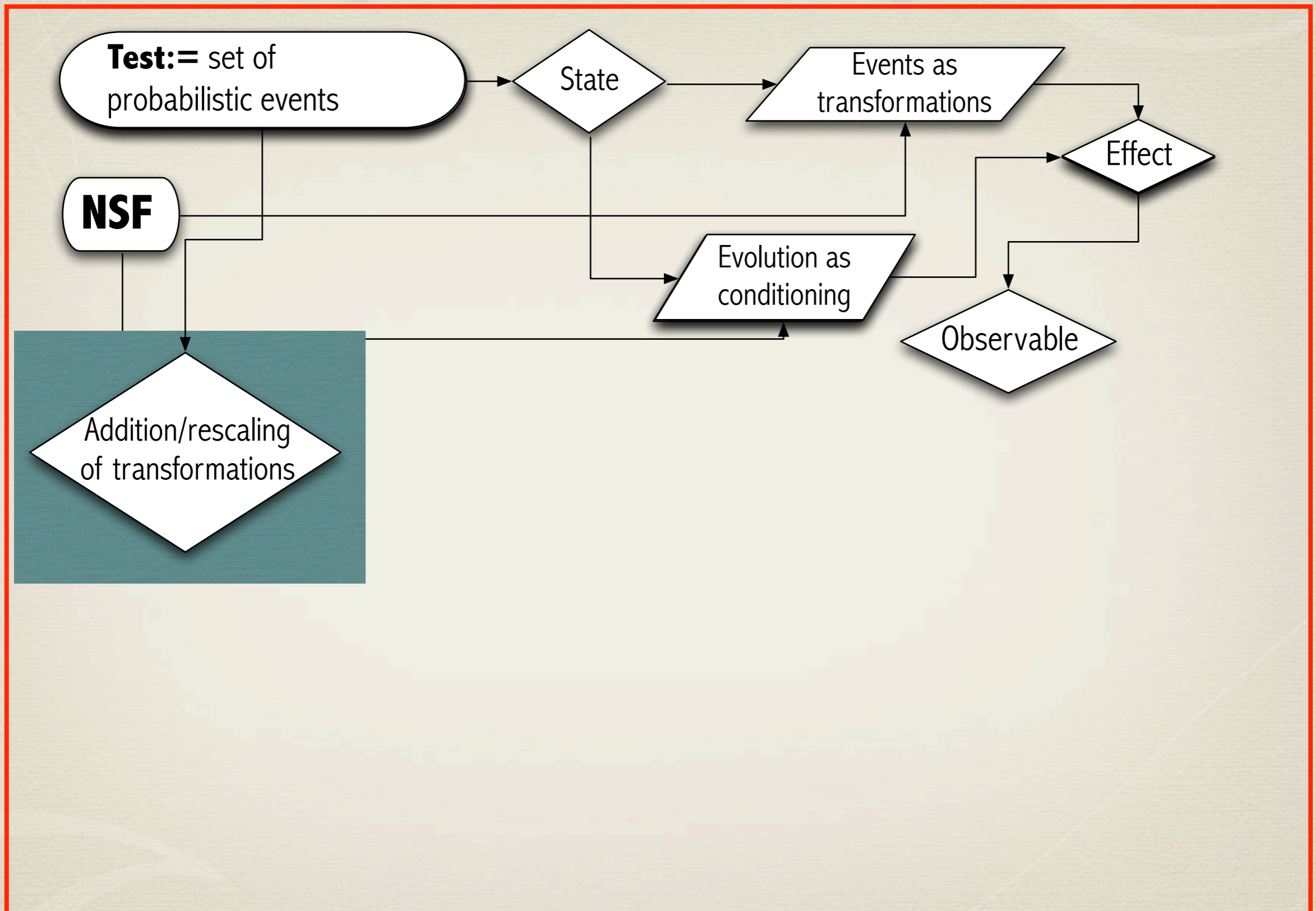
Logical flow-diagram



Logical flow-diagram



Logical flow-diagram



Addition of transformations

Two transformations \mathcal{A} and \mathcal{B} generally occurring in different tests are **test-compatible** if for every state ω one has

$$\omega(\mathcal{A}) + \omega(\mathcal{B}) \leq 1$$

For any two test-compatible transformations \mathcal{A}_1 and \mathcal{A}_2 we define the transformation $\mathcal{A}_1 + \mathcal{A}_2$ as the union event $\mathcal{A}_1 \cup \mathcal{A}_2$ as if they belong to the same test

$$(\mathcal{A}_1 + \mathcal{A}_2)\omega = \mathcal{A}_1\omega + \mathcal{A}_2\omega$$

Addition of transformations

Two transformations \mathcal{A} and \mathcal{B} generally occurring in different tests are **test-compatible** if for every state ω one has

$$\omega(\mathcal{A}) + \omega(\mathcal{B}) \leq 1$$

For any two test-compatible transformations \mathcal{A}_1 and \mathcal{A}_2 we define their sum $\mathcal{A}_1 + \mathcal{A}_2$ as follows:

$$\omega(\mathcal{A}_1 + \mathcal{A}_2) = \omega(\mathcal{A}_1) + \omega(\mathcal{A}_2) \quad (\text{probabilistic class})$$

$$\omega_{\mathcal{A}_1 + \mathcal{A}_2} = \frac{\omega(\mathcal{A}_1)}{\omega(\mathcal{A}_1 + \mathcal{A}_2)} \omega_{\mathcal{A}_1} + \frac{\omega(\mathcal{A}_2)}{\omega(\mathcal{A}_1 + \mathcal{A}_2)} \omega_{\mathcal{A}_2} \quad (\text{conditioning class})$$

Addition of transformations

Two transformations \mathcal{A} and \mathcal{B} generally occurring in different tests are **test-compatible** if for every state ω one has

$$\omega(\mathcal{A}) + \omega(\mathcal{B}) \leq 1$$

For any two test-compatible transformations \mathcal{A}_1 and \mathcal{A}_2 we define the transformation $\mathcal{A}_1 + \mathcal{A}_2$ as the union event $\mathcal{A}_1 \cup \mathcal{A}_2$ as if they belong to the same test

$$(\mathcal{A}_1 + \mathcal{A}_2)\omega = \mathcal{A}_1\omega + \mathcal{A}_2\omega$$

Rescaling of transformations

The **rescaled transformation** $\lambda\mathcal{A}$ of \mathcal{A} , $\lambda \in [0, 1]$ is the transformation giving the same conditioning but occurring with probability rescaled by λ for all states.

Rescaling of transformations

The **rescaled transformation** $\lambda\mathcal{A}$ of \mathcal{A} , $\lambda \in [0, 1]$ is the transformation giving the same conditioning but occurring with probability rescaled by λ for all states.

Atomic: a transformation that cannot be “nontrivially” refined in any test, i.e. it cannot be written as $\mathcal{A} = \sum_i \mathcal{A}_i$ with $\mathcal{A}_i \neq \lambda_i \mathcal{A}$ for some i and $0 < \lambda_i < 1$.

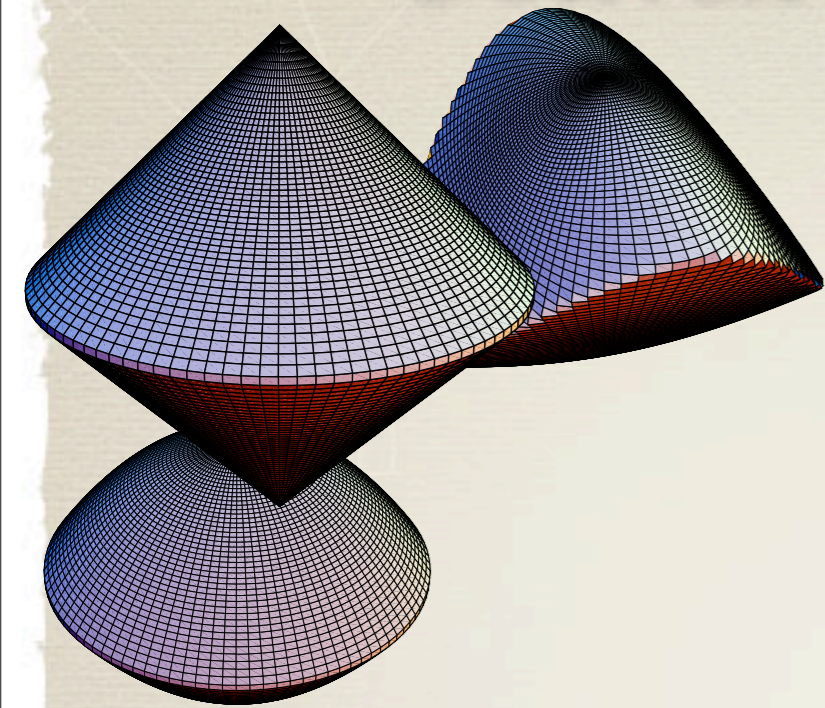
Rescaling of transformations

The **rescaled transformation** $\lambda\mathcal{A}$ of \mathcal{A} , $\lambda \in [0, 1]$ is the transformation giving the same conditioning but occurring with probability rescaled by λ for all states.

Atomic: a transformation that cannot be “nontrivially” refined in any test, i.e. it cannot be written as $\mathcal{A} = \sum_i \mathcal{A}_i$ with $\mathcal{A}_i \neq \lambda_i \mathcal{A}$ for some i and $0 < \lambda_i < 1$.

[Notice: the identity transformation \mathcal{I} is not necessarily atomic]

Convex sets, Cones and Linear spaces



Convex set of states:

\mathcal{S} , cone: \mathcal{S}_+

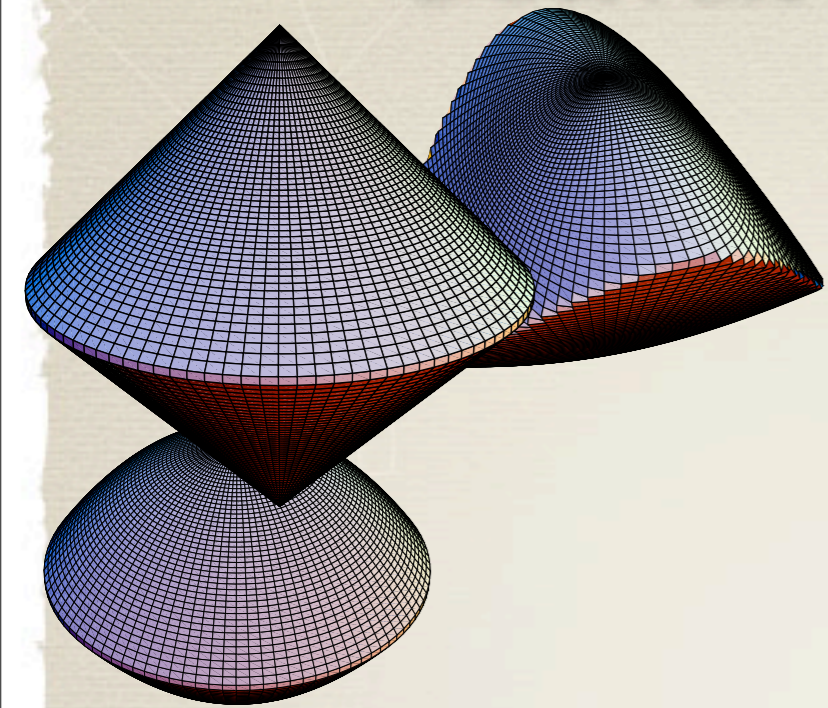
Convex set of effects:

\mathcal{E} , cone: \mathcal{E}_+

Convex monoid of transformations:

\mathcal{T} , cone: \mathcal{T}_+

Convex sets, Cones and Linear spaces



Convex set of states:

\mathcal{S} , cone: \mathcal{S}_+

Convex set of effects:

\mathcal{E} , cone: \mathcal{E}_+

Convex monoid of transformations:

\mathcal{T} , cone: \mathcal{T}_+

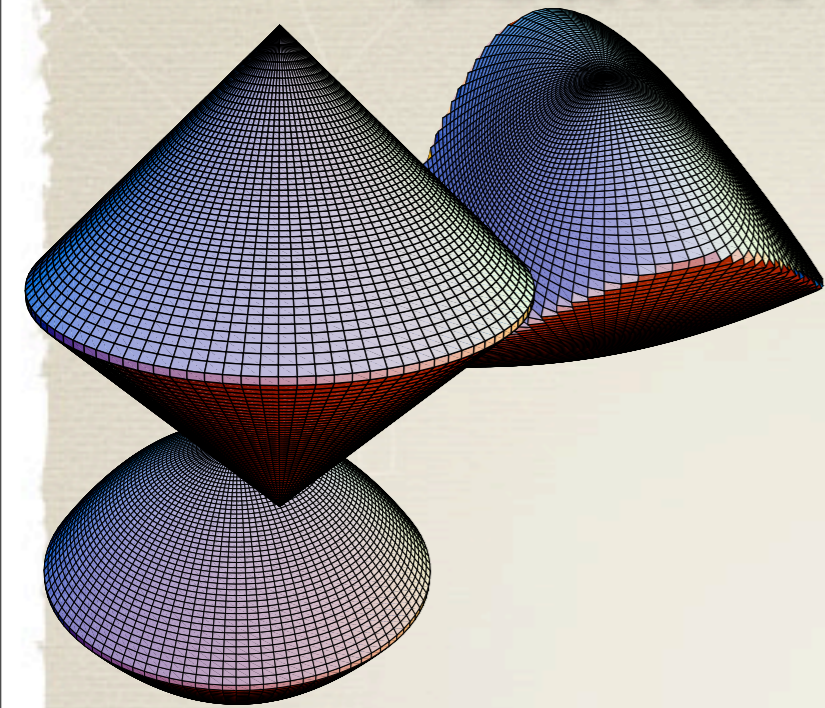
Linear spaces:

$$\mathcal{S}_{\mathbb{R}} = \text{Span}_{\mathbb{R}} \mathcal{S}$$

$$\mathcal{S}_{\mathbb{C}} = \text{Span}_{\mathbb{C}} \mathcal{S}$$

$$\mathcal{E}_{\mathbb{R}}, \mathcal{E}_{\mathbb{C}}, \mathcal{T}_{\mathbb{R}}, \mathcal{T}_{\mathbb{C}}$$

Convex sets, Cones and Linear spaces



Convex set of states:

$$\mathcal{S}, \text{ cone: } \mathcal{S}_+$$

Convex set of effects:

$$\mathcal{E}, \text{ cone: } \mathcal{E}_+$$

Convex monoid of transformations:

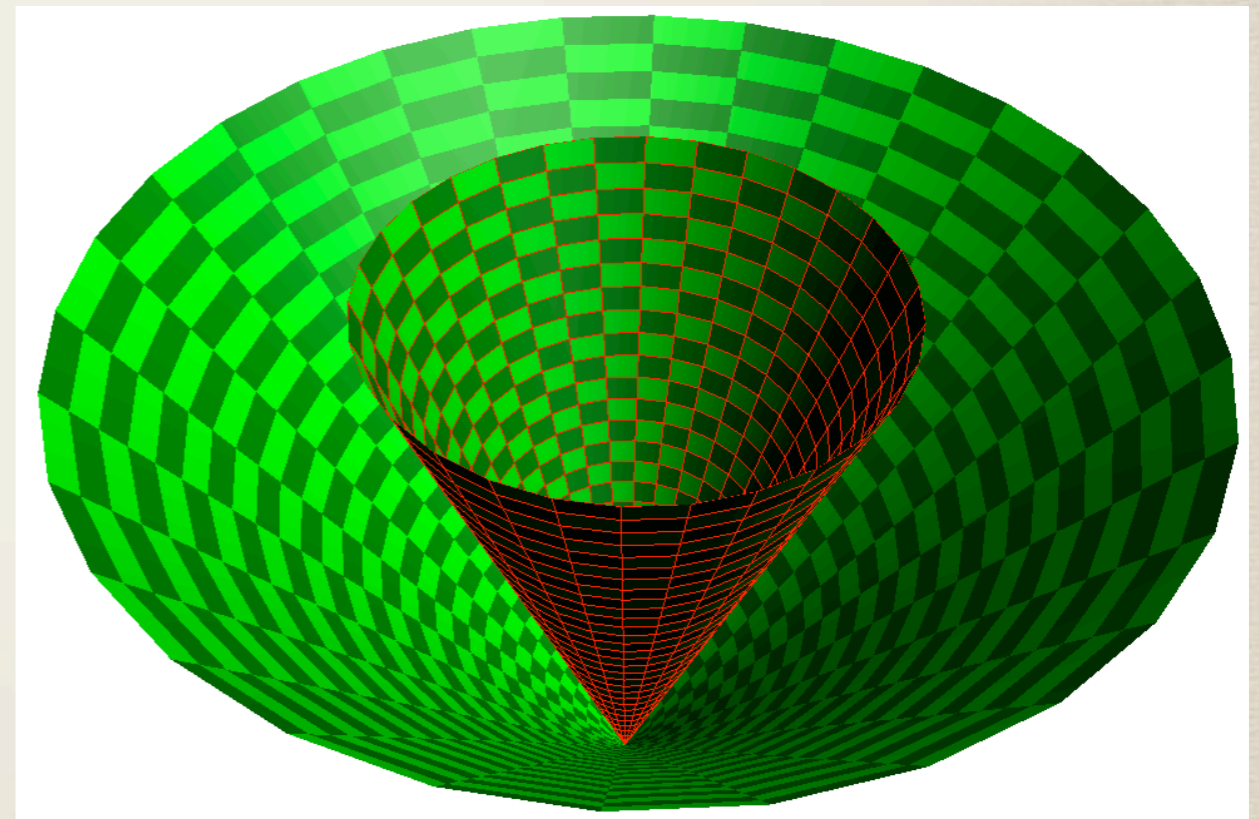
$$\mathcal{T}, \text{ cone: } \mathcal{T}_+$$

Linear spaces:

$$\mathcal{S}_{\mathbb{R}} = \text{Span}_{\mathbb{R}} \mathcal{S}$$

$$\mathcal{S}_{\mathbb{C}} = \text{Span}_{\mathbb{C}} \mathcal{S}$$

$$\mathcal{E}_{\mathbb{R}}, \mathcal{E}_{\mathbb{C}}, \mathcal{T}_{\mathbb{R}}, \mathcal{T}_{\mathbb{C}}$$



Hypothesis of no limitation to preparability: $\mathcal{S}_+ = (\mathcal{E}_+)^*$

Informational completeness

Informational completeness

Informationally complete observable: \mathbb{L}

$$\mathcal{E}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{L})$$

Informational completeness

Informationally complete observable: \mathbb{L}

$$\mathcal{E}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{L})$$

Separating set of states: \mathbb{S}

$$\mathcal{G}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{S})$$

Informational completeness

Informationally complete observable: \mathbb{L}

$$\mathcal{E}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{L})$$

Separating set of states: \mathbb{S}

$$\mathcal{G}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{S})$$

Quantum Bureau International des Poids et Mesures (Fuchs):

$$\mathbb{S} = \{\mathcal{I}_i\}$$

$\mathcal{I}_i \omega = \omega(\mathcal{I}_i) \omega_i$, $\forall \omega \in \mathcal{G}$, $\{\omega_i\}$ separating

$\{[\mathcal{I}_i]_{\text{eff}}\}$ informationally complete observable

Informational completeness

Informationally complete observable: \mathbb{L}

$$\mathcal{E}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{L})$$

Separating set of states: \mathbb{S}

$$\mathcal{G}_{\mathbb{R}} = \text{Span}_{\mathbb{R}}(\mathbb{S})$$

C^* -algebra of transformations (finite dim.)

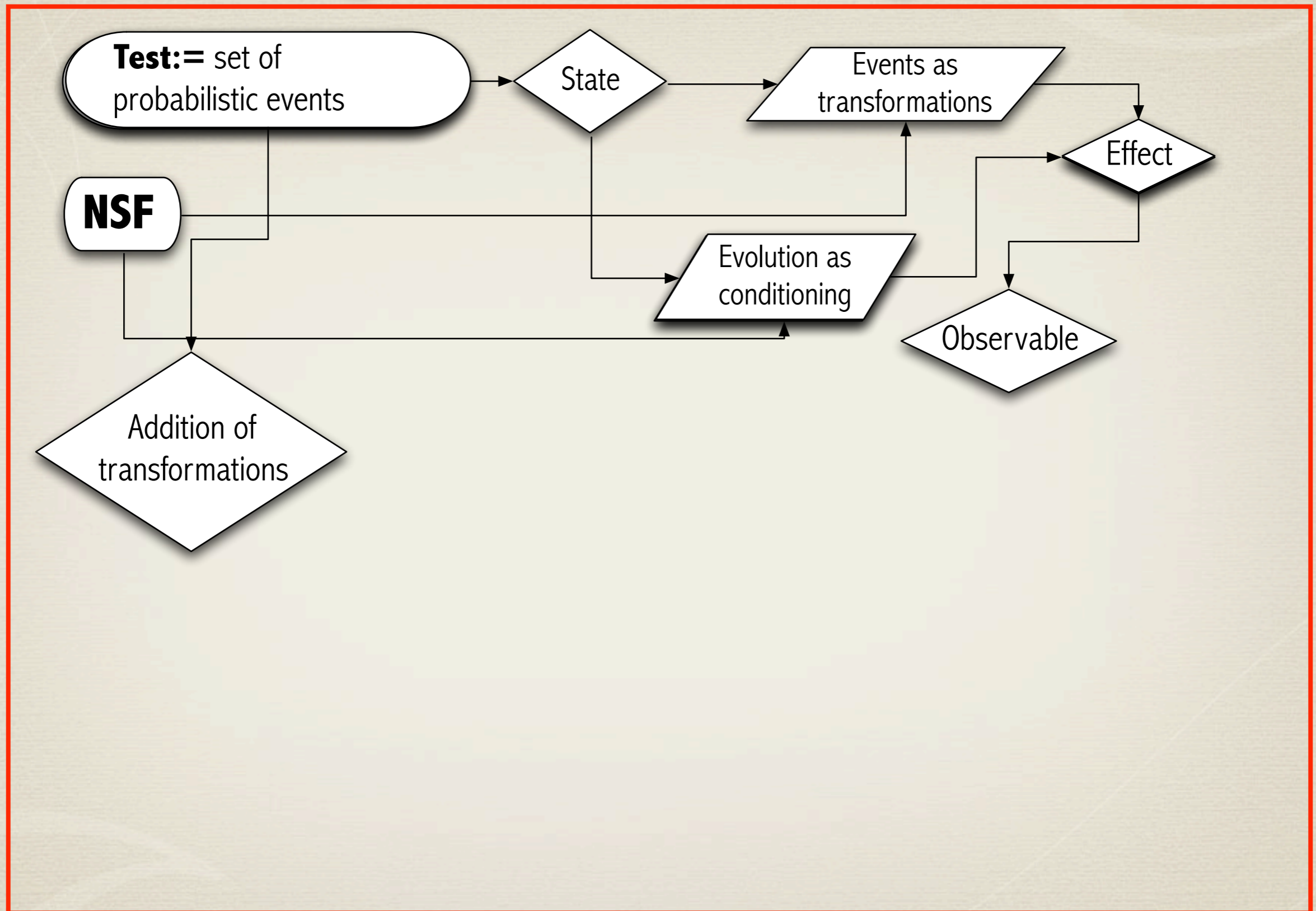
Transformations/events are linear maps over effects, i.e. they make a **matrix algebra** over effects (or over states)

C^* -algebra of transformations (finite dim.)

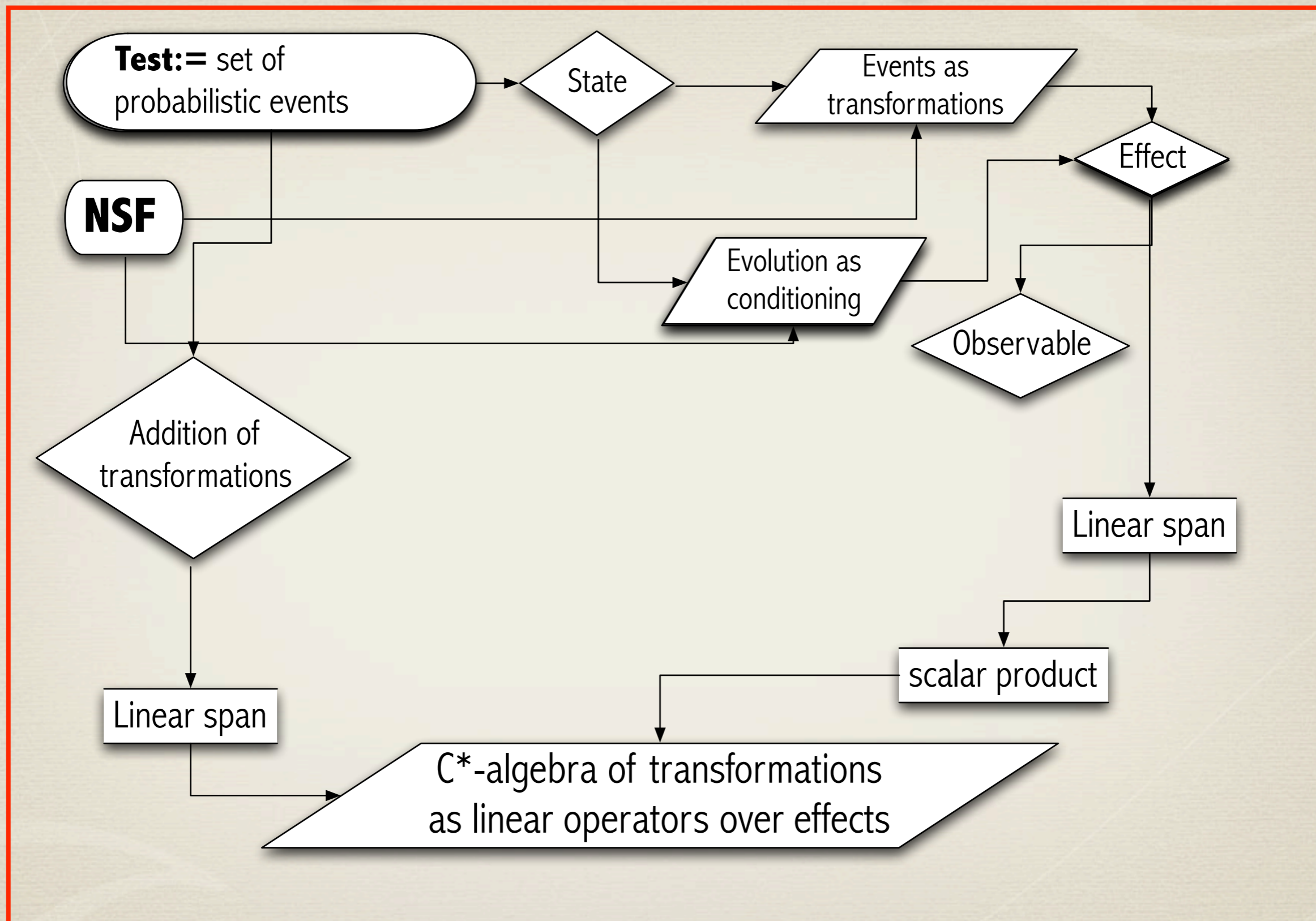
Transformations/events are linear maps over effects, i.e. they make a **matrix algebra** over effects (or over states)

One can introduce a scalar product over effects ...
 \Rightarrow transformations become a C^* -algebra ...

Logical flow-diagram

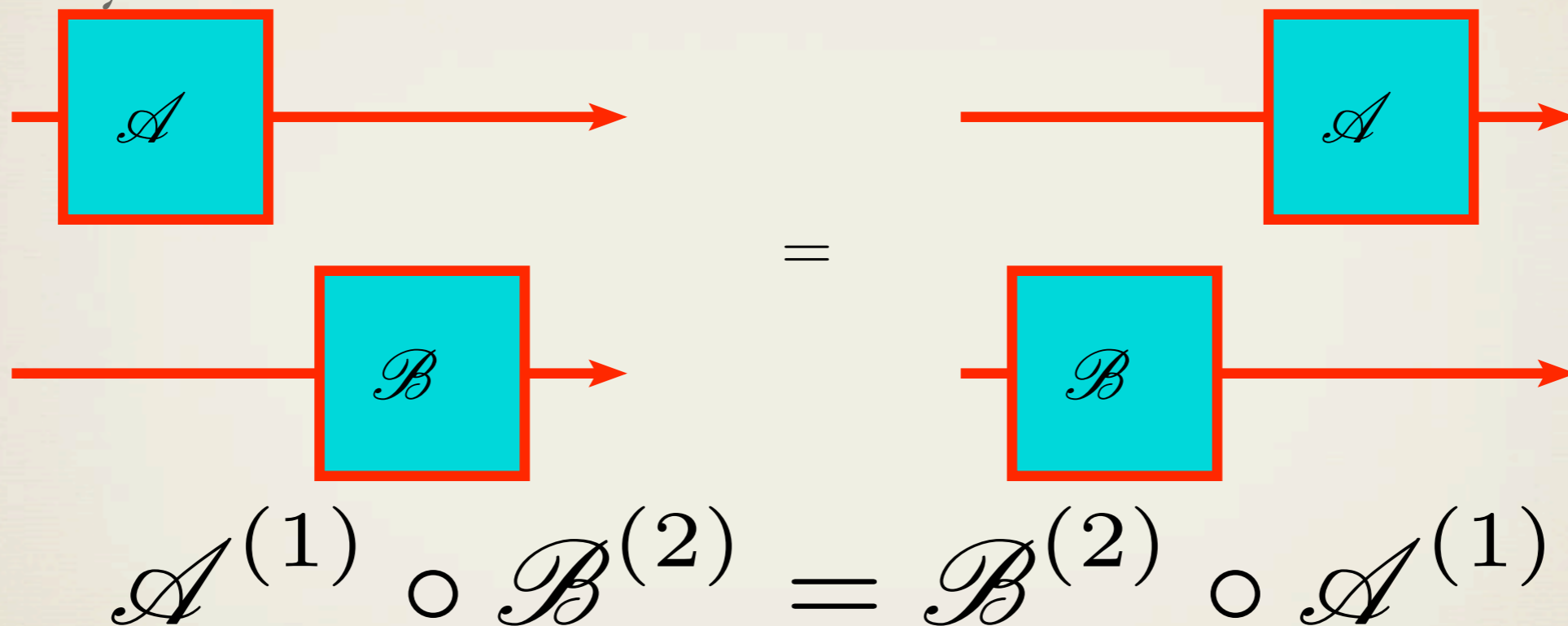


Logical flow-diagram



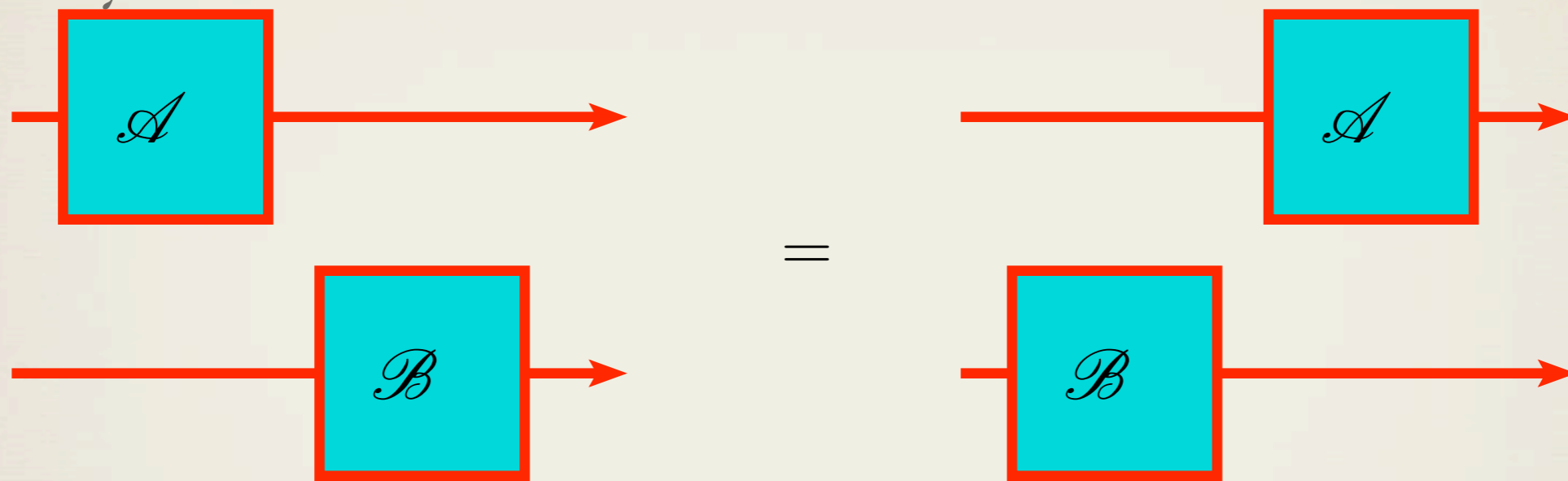
INDEPENDENT SYSTEMS

Two systems are **independent** if on each system it is possible to perform all their tests as **local tests**, i.e. such that on every joint state one has the commutativity of the transformations from different systems



INDEPENDENT SYSTEMS

Two systems are **independent** if on each system it is possible to perform all their tests as **local tests**, i.e. such that on every joint state one has the commutativity of the transformations from different systems



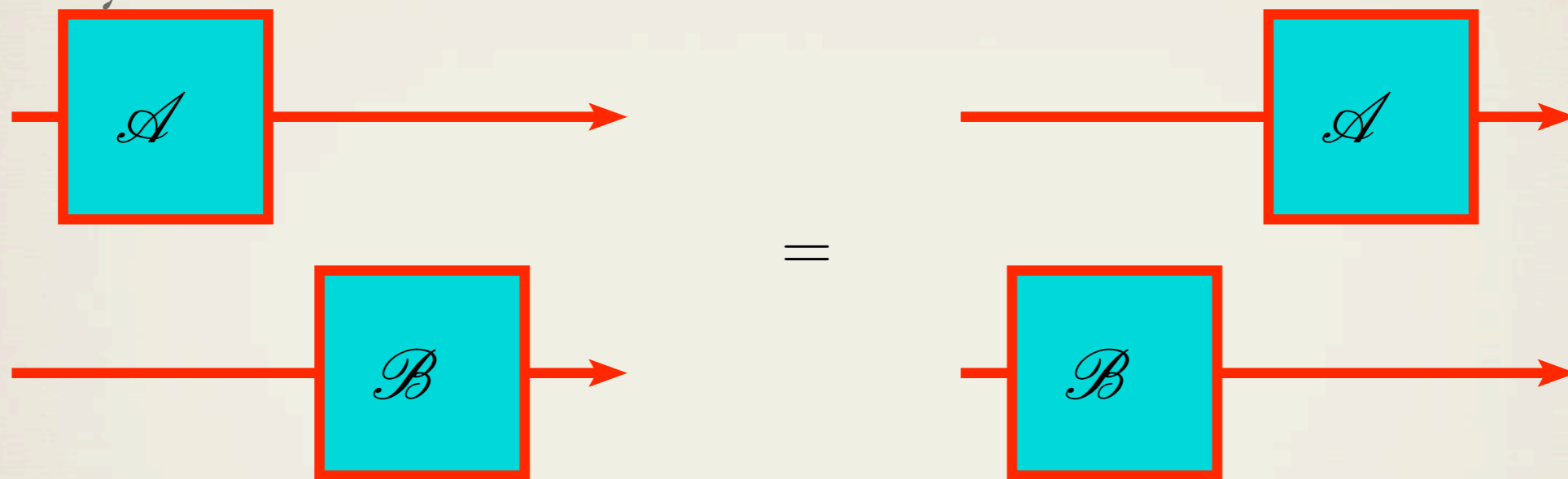
$$A^{(1)} \circ B^{(2)} = B^{(2)} \circ A^{(1)}$$

$$(A, B, C, \dots) \doteq A^{(1)} \circ B^{(2)} \circ C^{(3)} \circ \dots$$

$$S_1 \times S_2$$

INDEPENDENT SYSTEMS

Two systems are **independent** if on each system it is possible to perform all their tests as **local tests**, i.e. such that on every joint state one has the commutativity of the transformations from different systems



$$A^{(1)} \circ B^{(2)} = B^{(2)} \circ A^{(1)}$$

$$(A, B, C, \dots) \doteq A^{(1)} \circ B^{(2)} \circ C^{(3)} \circ \dots$$

$$S_1 \times S_2$$

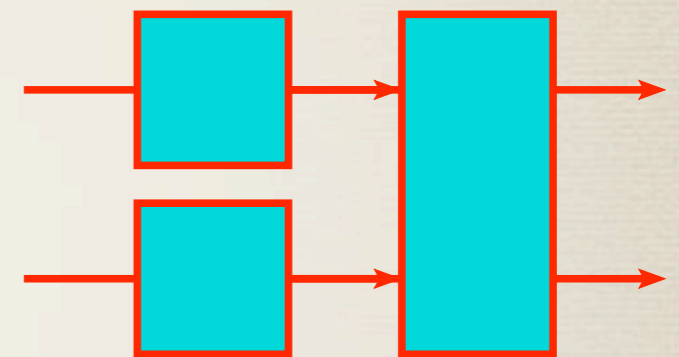
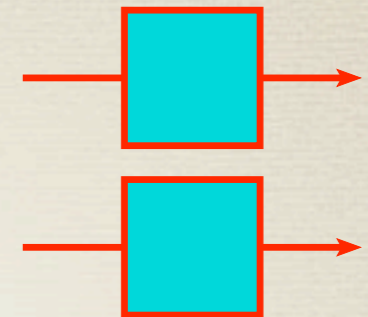
$$[(A, B, C, \dots)]_{\text{eff}} = ([A]_{\text{eff}}, [B]_{\text{eff}}, [C]_{\text{eff}}, \dots)$$

COMPOSTING SYSTEMS

We compose the two systems S_1 and S_2 into the bipartite system $S_1 \odot S_2$ considered as a new system containing all **local tests** $S_1 \times S_2$ plus other tests, and closing w.r.t. coarse graining, convex combination and cascading:

$$S_1 \odot S_2 \supseteq S_1 \times S_2$$

Nonlocal tests: $S_1 \odot S_2 \setminus S_1 \times S_2$



MARGINAL STATE

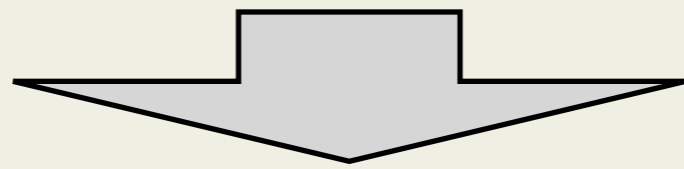
For a multipartite system we define the marginal state $\Omega|_n$ of the n -th system the state that gives the probability of any local transformation \mathcal{A} on the n -th system with all other systems untouched, namely

$$\Omega|_n(\mathcal{A}) := \Omega(\mathcal{I}, \dots, \mathcal{I}, \underbrace{\mathcal{A}}_{n\text{-th}}, \mathcal{I}, \dots)$$

MARGINAL STATE

For a multipartite system we define the marginal state $\Omega|_n$ of the n -th system the state that gives the probability of any local transformation \mathcal{A} on the n -th system with all other systems untouched, namely

$$\Omega|_n(\mathcal{A}) := \Omega(\mathcal{I}, \dots, \mathcal{I}, \underbrace{\mathcal{A}}_{n\text{-th}}, \mathcal{I}, \dots)$$

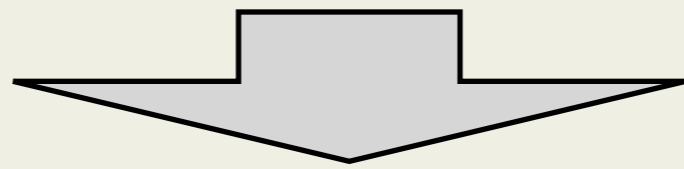


$$\Omega|_n(a) \doteq \Omega(e, \dots, e, \underbrace{a}_{n\text{th}}, e, \dots)$$

MARGINAL STATE

For a multipartite system we define the marginal state $\Omega|_n$ of the n -th system the state that gives the probability of any local transformation \mathcal{A} on the n -th system with all other systems untouched, namely

$$\Omega|_n(\mathcal{A}) := \Omega(\mathcal{I}, \dots, \mathcal{I}, \underbrace{\mathcal{A}}_{n\text{-th}}, \mathcal{I}, \dots)$$



$$\Omega|_n(a) \doteq \Omega(e, \dots, e, \underbrace{a}_{n\text{th}}, e, \dots)$$

NS: (no-signaling) any local test on a system is equivalent to no-test on another independent system.



PROBABILISTIC
THEORY?

PROBABILISTIC THEORY?

Matrix algebra of
transformations over effects!

PROBABILISTIC THEORY?

Matrix algebra of
transformations over effects!

Independent
systems =
no-signaling

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: \mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: $\mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+$

$$\mathcal{S} = \{\zeta, \omega, \dots, A, B, C, \dots, a, b, \dots\}$$



System



tests

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: \mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+

$$\mathcal{S} = \{\zeta, \omega, \dots, A, B, C, \dots, a, b, \dots\}$$



System



tests

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$



variable

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: \mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+

$$\mathcal{S} = \{\zeta, \omega, \dots, \mathbb{A}, \mathbb{B}, \mathbb{C}, \dots, a, b, \dots\}$$



System



tests

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$

variable

$$\mathcal{A} \in a \leftarrow \text{effect}$$

transformation

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: \mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+

$$\mathcal{S} = \{\zeta, \omega, \dots, A, B, C, \dots, a, b, \dots\}$$



System



tests

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$

variable

$$\mathcal{A} \in a \leftarrow \text{effect}$$

transformation

e deterministic effect

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: $\mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+$

$$\mathcal{S} = \{\zeta, \omega, \dots, A, B, C, \dots, a, b, \dots\}$$



System



tests

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$

variable

$$\mathcal{A} \in a \leftarrow \text{effect}$$

transformation

Bipartite: $\mathcal{E}(\mathcal{S}_1 \odot \mathcal{S}_2)$

$$\mathcal{E}^{\odot 2} := \mathcal{E}(\mathcal{S}^{\odot 2})$$

e deterministic effect

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: $\mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+$

$$\mathcal{S} = \{\zeta, \omega, \dots, \mathbb{A}, \mathbb{B}, \mathbb{C}, \dots, a, b, \dots\}$$



System



tests

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$

variable

$$\mathcal{A} \in a \leftarrow \text{effect}$$

transformation

Bipartite: $\mathcal{E}(\mathcal{S}_1 \odot \mathcal{S}_2)$

$$\mathcal{E}^{\odot 2} := \mathcal{E}(\mathcal{S}^{\odot 2})$$

e deterministic effect

$$\omega, \sigma \in \mathcal{S}$$

$$a, b \in \mathcal{E}$$

$$\Omega, \Phi \in \mathcal{S}^{\odot 2}$$

$$E \in \mathcal{E}^{\odot 2}$$

Review of notation

Convex sets:

\mathcal{S} states

\mathcal{E} effects

\mathcal{T} transformations

cones: $\mathcal{S}_+ \mathcal{E}_+ \mathcal{T}_+$

$$\mathcal{S} = \{\zeta, \omega, \dots, \mathbb{A}, \mathbb{B}, \mathbb{C}, \dots, a, b, \dots\}$$



System



tests

$$\mathcal{A}\omega = \omega(\cdot \circ \mathcal{A})$$

variable

$$\mathcal{A} \in a \leftarrow \text{effect}$$

transformation

Bipartite: $\mathcal{E}(\mathcal{S}_1 \odot \mathcal{S}_2)$

$$\mathcal{E}^{\odot 2} := \mathcal{E}(\mathcal{S}^{\odot 2})$$

e deterministic effect

$$\omega, \sigma \in \mathcal{S}$$

$$a, b \in \mathcal{E}$$

$$\Omega, \Phi \in \mathcal{S}^{\odot 2}$$

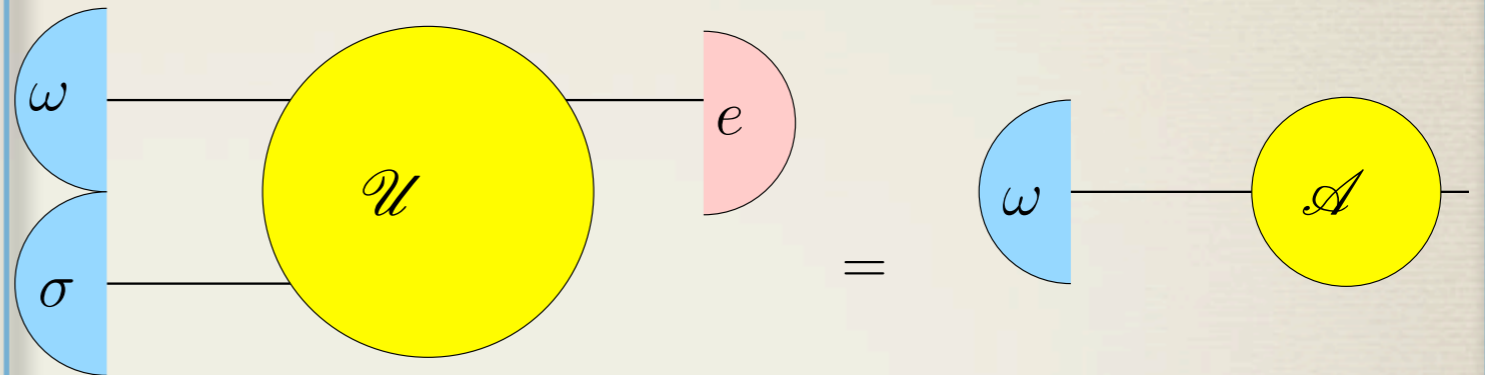
$$E \in \mathcal{E}^{\odot 2}$$

$$\mathcal{U}_{12}, (\mathcal{A}, \mathcal{I}) \in \mathcal{T}^{\odot 2}$$

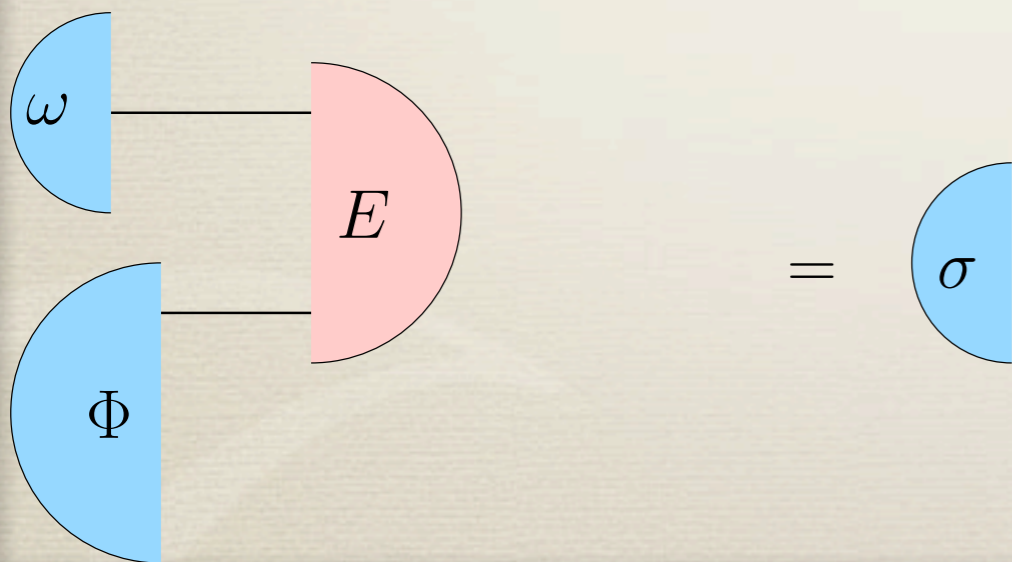
Review of notation

$$\omega(a) \equiv a(\omega)$$

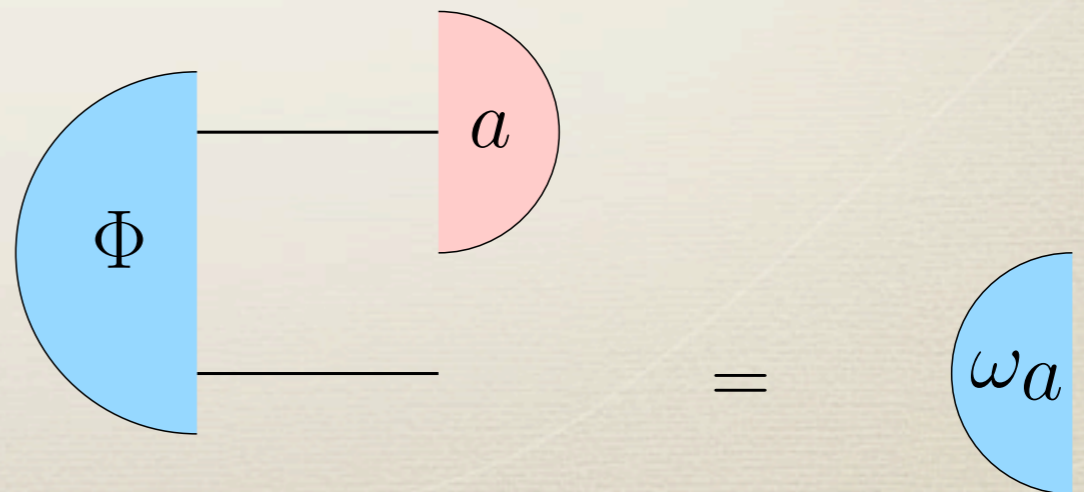
$$\mathcal{U}_{12}(\sigma, \omega)(e, \cdot) = \mathcal{A}\omega$$



$$E_{23}(\Phi, \omega) = \sigma \in \mathfrak{S}$$



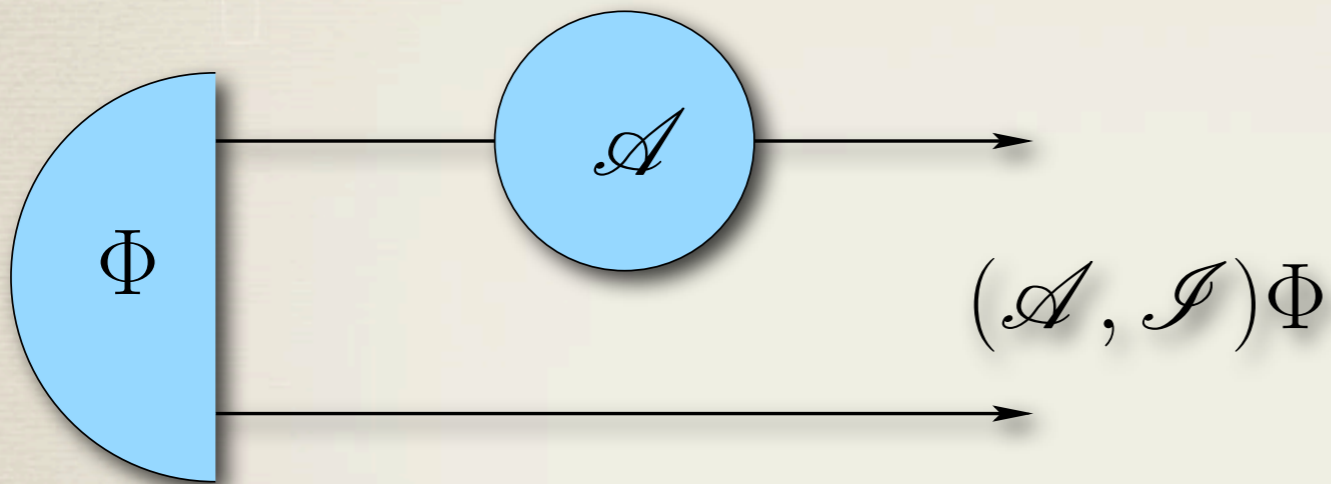
$$\Phi(a, \cdot) = \omega_a \in \mathfrak{S}_+$$



FAITHFUL STATES

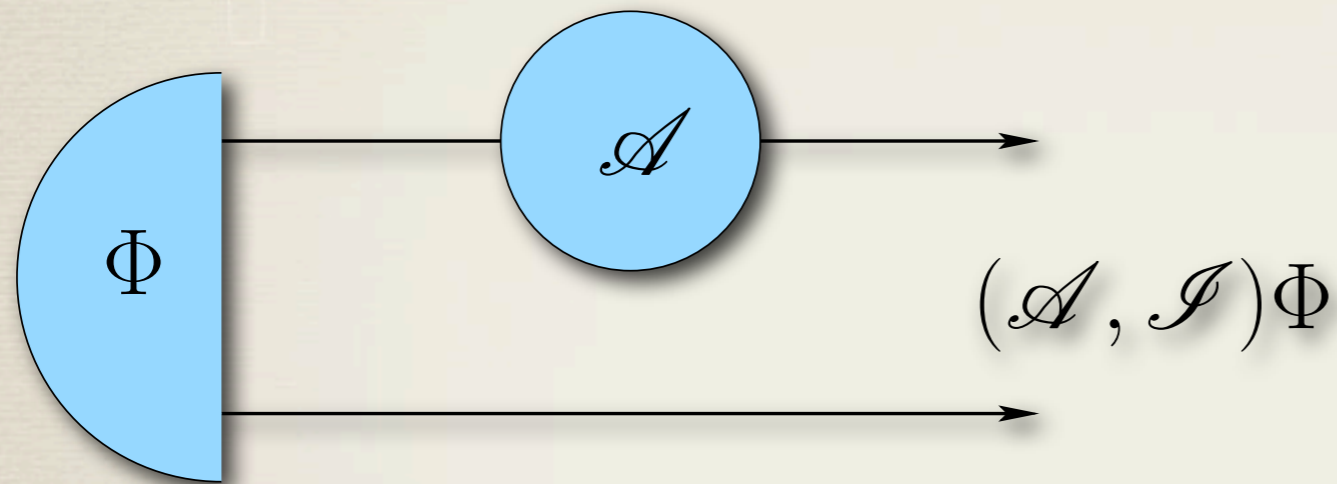
FAITHFUL STATES

A state Φ of a bipartite system is **dynamically faithful** when the output state $(\mathcal{A}, \mathcal{I})\Phi$ from a local transformation \mathcal{A} on one system is in 1-to-1 correspondence with the transformation \mathcal{A}



FAITHFUL STATES

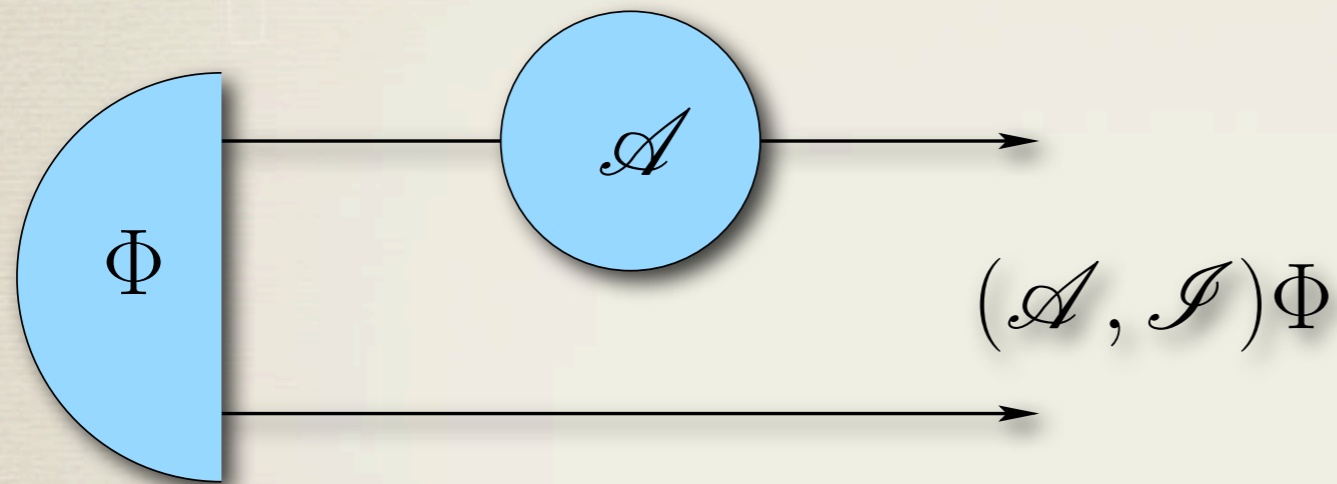
A state Φ of a bipartite system is **dynamically faithful** when the output state $(\mathcal{A}, \mathcal{I})\Phi$ from a local transformation \mathcal{A} on one system is in 1-to-1 correspondence with the transformation \mathcal{A}



calibrability of tests

FAITHFUL STATES

A state Φ of a bipartite system is **dynamically faithful** when the output state $(\mathcal{A}, \mathcal{I})\Phi$ from a local transformation \mathcal{A} on one system is in 1-to-1 correspondence with the transformation \mathcal{A}



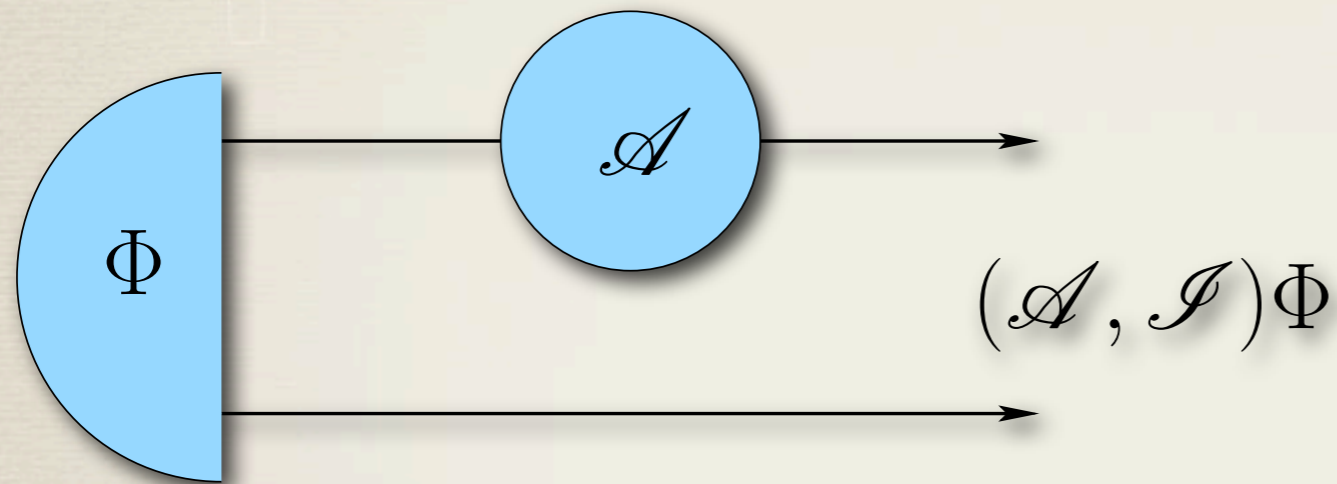
calibrability of tests

A state Φ of a bipartite system is **preparationally faithful** if every joint state Ψ can be achieved by a suitable local transformation \mathcal{T}_Ψ on one system occurring with nonzero probability



FAITHFUL STATES

A state Φ of a bipartite system is **dynamically faithful** when the output state $(\mathcal{A}, \mathcal{I})\Phi$ from a local transformation \mathcal{A} on one system is in 1-to-1 correspondence with the transformation \mathcal{A}



calibrability of tests

local state-preparability

A state Φ of a bipartite system is **preparationally faithful** if every joint state Ψ can be achieved by a suitable local transformation \mathcal{T}_Ψ on one system occurring with nonzero probability



Postulate PFAITH

PFAITH: For any couple of identical systems, there exist a symmetric* state Φ that is preparationally faithful.

(*) invariant under permutation of the two systems

Postulate PFAITH

PFAITH: For any couple of identical systems, there exist a symmetric* state Φ that is preparationally faithful.

Theorem: Φ is also dynamically faithful.

(*) invariant under permutation of the two systems

Consequences of PFAITH

Consequences of PFAITH

- ▶ Calibrability & Preparability by just a single preparation

Consequences of PFAITH

- ▶ Calibrability & Preparability by just a single preparation
- ▶ Impossibility of secure bit commitment

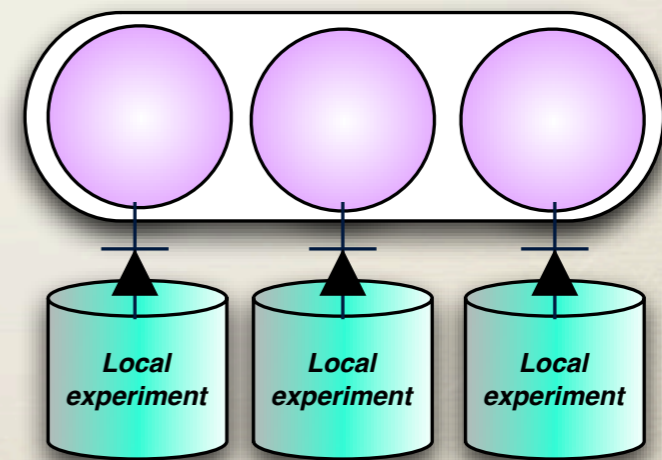
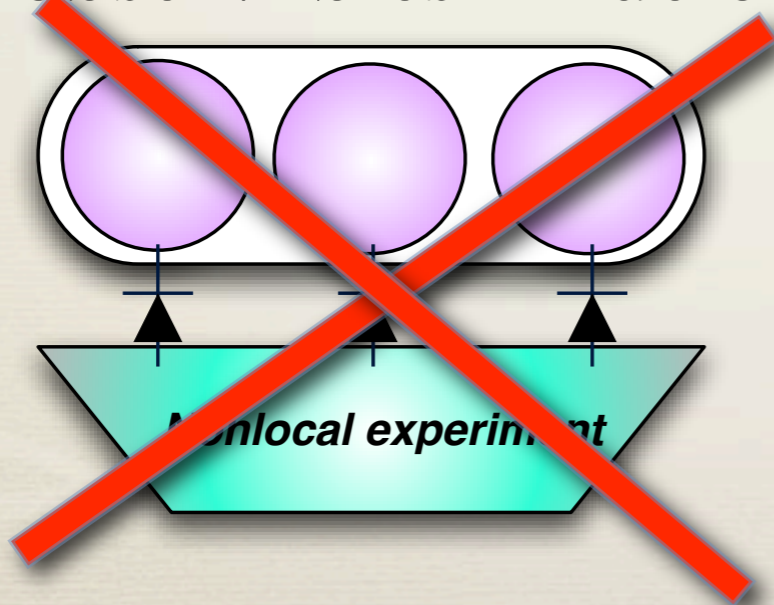
Consequences of PFAITH

- ▶ Calibrability & Preparability by just a single preparation
- ▶ Impossibility of secure bit commitment
- ▶ Marginal state $\chi = \Phi(e, \cdot)$ internal and invariant under a “transposed” deterministic test

Consequences of PFAITH

- ▶ Calibrability & Preparability by just a single preparation
- ▶ Impossibility of secure bit commitment
- ▶ Marginal state $\chi = \Phi(e, \cdot)$ internal and invariant under a “transposed” deterministic test

- ▶ **Local observability:** There exist global info-complete observables made of local info-complete



Holism



Reductionism

Consequences of PFAITH

$$\mathbb{F} = \mathbb{R}, \mathbb{C}$$

Consequences of PFAITH

$$\blacktriangleright \mathfrak{S}_{\mathbb{F}}(S^{\odot 2}) \simeq \mathfrak{I}_{\mathbb{F}}(S)$$

$$\mathbb{F} = \mathbb{R}, \mathbb{C}$$

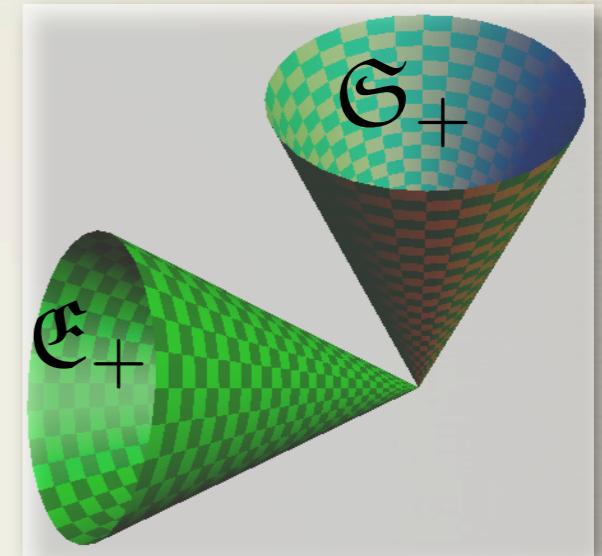
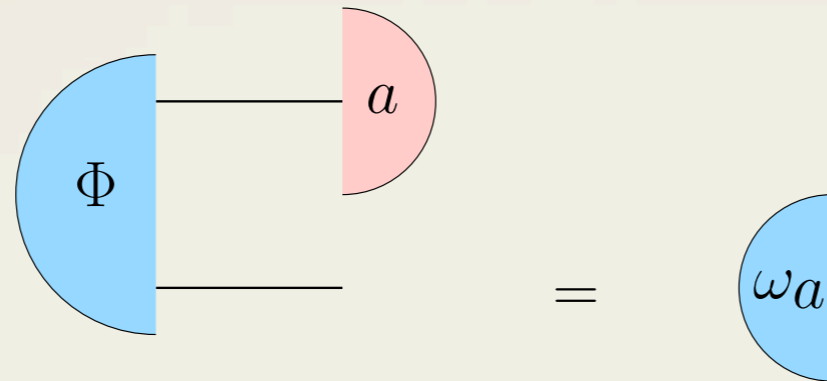
Consequences of PFAITH

$$\triangleright \mathfrak{S}_{\mathbb{F}}(\mathcal{S}^{\odot 2}) \simeq \mathfrak{T}_{\mathbb{F}}(\mathcal{S})$$

$$\mathbb{F} = \mathbb{R}, \mathbb{C}$$

Weak self-duality: State and effect cones are isomorphic:

$$\mathfrak{E}_+ \ni a \mapsto \omega_a = \Phi(a, \cdot) \in \mathfrak{S}_+$$



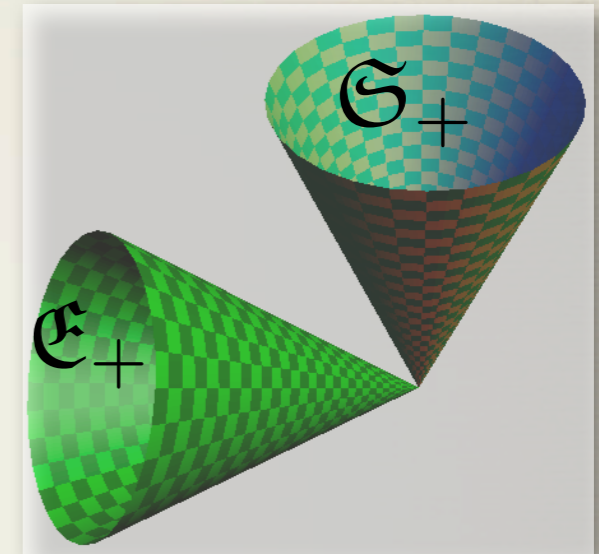
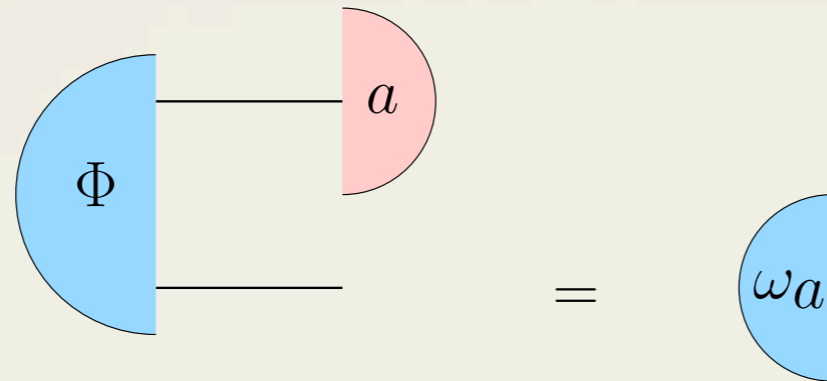
Consequences of PFAITH

$$\mathfrak{S}_{\mathbb{F}}(S^{\odot 2}) \simeq \mathfrak{I}_{\mathbb{F}}(S)$$

$$\mathbb{F} = \mathbb{R}, \mathbb{C}$$

▶ **Weak self-duality:** State and effect cones are isomorphic:

$$\mathfrak{E}_+ \ni a \mapsto \omega_a = \Phi(a, \cdot) \in \mathfrak{S}_+$$



▶ **Tensor product representation:**

$$\mathfrak{E}_{\mathbb{F}}(S^{\odot 2}) = \mathfrak{E}_{\mathbb{F}}(S)^{\otimes 2}$$

$$\mathfrak{S}_{\mathbb{F}}(S^{\odot 2}) = \mathfrak{S}_{\mathbb{F}}(S)^{\otimes 2}$$

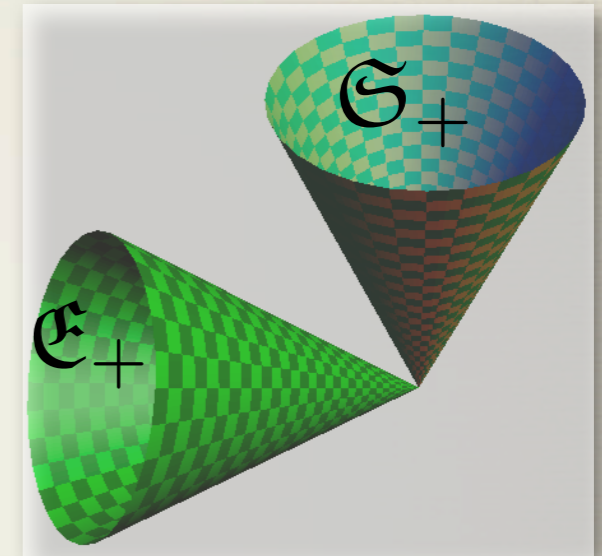
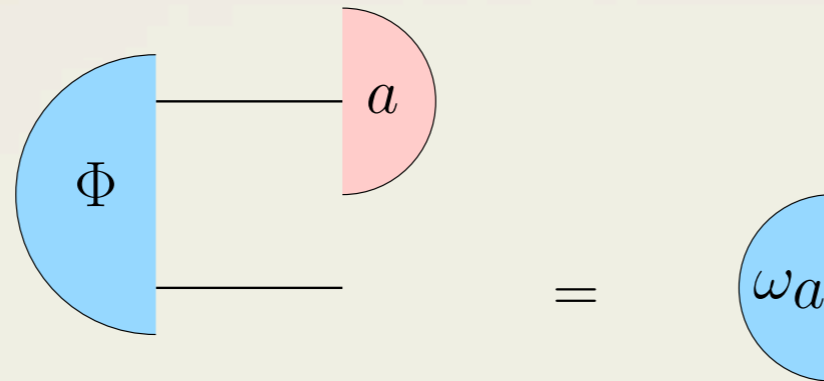
Consequences of PFAITH

$$\mathfrak{S}_{\mathbb{F}}(S^{\odot 2}) \simeq \mathfrak{T}_{\mathbb{F}}(S)$$

$$\mathbb{F} = \mathbb{R}, \mathbb{C}$$

► **Weak self-duality:** State and effect cones are isomorphic:

$$\mathfrak{E}_+ \ni a \mapsto \omega_a = \Phi(a, \cdot) \in \mathfrak{S}_+$$



► **Tensor product representation:**

$$\mathfrak{E}_{\mathbb{F}}(S^{\odot 2}) = \mathfrak{E}_{\mathbb{F}}(S)^{\otimes 2}$$

$$\mathfrak{S}_{\mathbb{F}}(S^{\odot 2}) = \mathfrak{S}_{\mathbb{F}}(S)^{\otimes 2}$$

► **Space of transformations is complete:**

$$\mathfrak{T}_{\mathbb{F}} = \text{Lin}(\mathfrak{E}_{\mathbb{F}})$$

Consequences of PFAITH

Consequences of PFAITH

The faithful state Φ provides a non-degenerate **scalar product** over effects via its Jordan form (ζ Jordan involution):

$$\forall a, b \in \mathcal{E}_{\mathbb{R}}, \quad \Phi(b|a)_{\Phi} := |\Phi|(b, a) = \Phi(\zeta(b), a)$$

Consequences of PFAITH

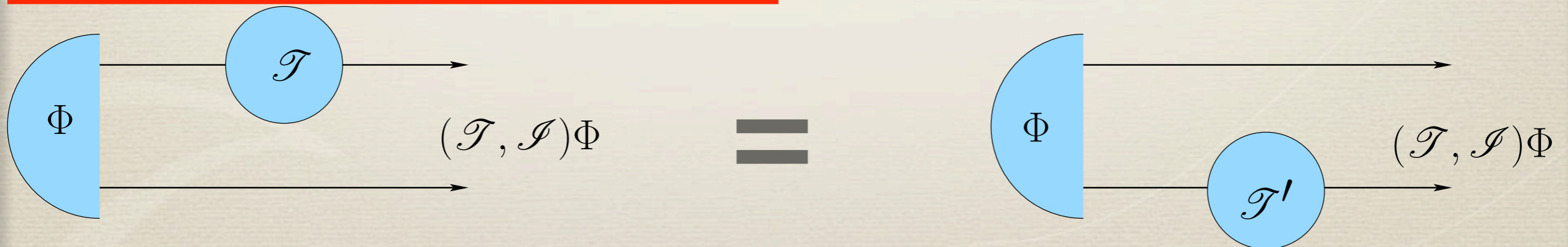
The faithful state Φ provides a non-degenerate **scalar product** over effects via its Jordan form (ζ Jordan involution):

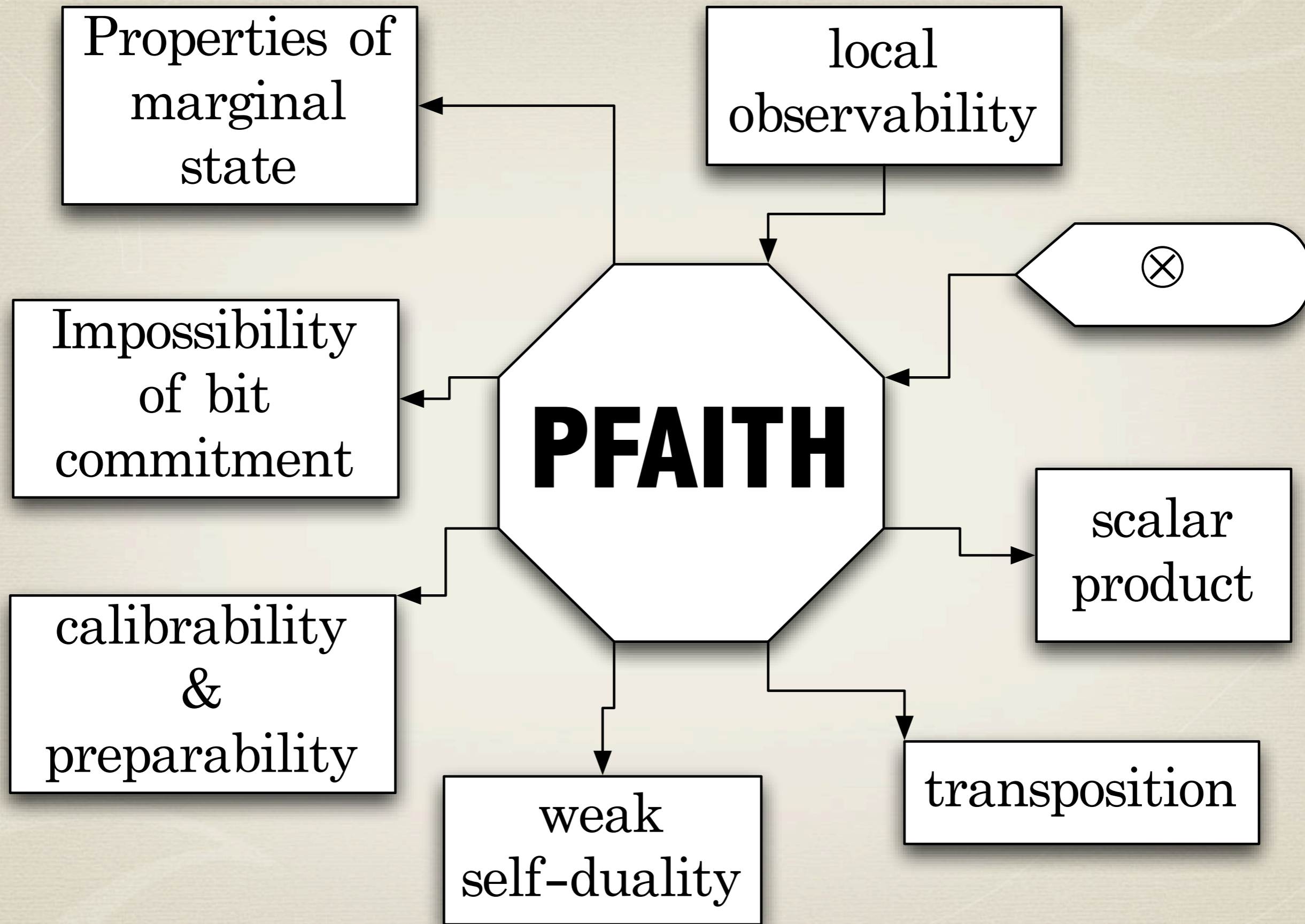
$$\forall a, b \in \mathcal{E}_{\mathbb{R}}, \quad \Phi(b|a)_{\Phi} := |\Phi|(b, a) = \Phi(\zeta(b), a)$$

It allows to introduce an **operational notion of transposition** for transformations:

$$(\mathcal{J}, \mathcal{I})\Phi = (\mathcal{I}, \mathcal{J}')\Phi$$

1. $(\mathcal{A} + \mathcal{B})' = \mathcal{A}' + \mathcal{B}'$
2. $(\mathcal{A}')' = \mathcal{A}$,
3. $(\mathcal{A} \circ \mathcal{B})' = \mathcal{B}' \circ \mathcal{A}'$



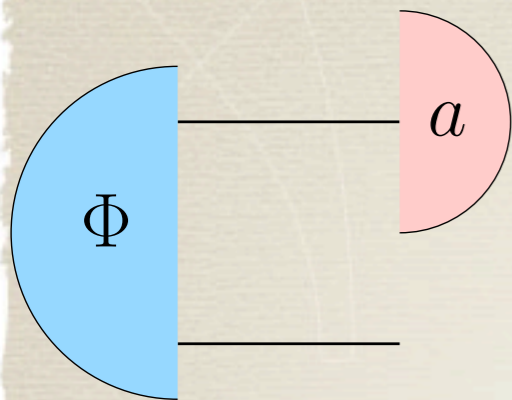


INTERLUDE

Exploring
Postulates:

FAITHE and
PURIFY

Faithful effect



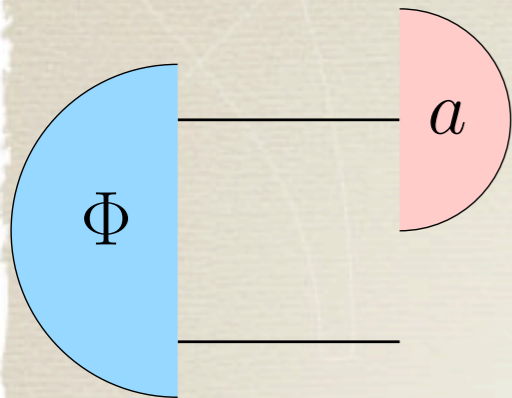
=



Remind the cone-isomorphism from the faithful state Φ

$$\mathfrak{E}_+ \ni a \mapsto \omega_a = \Phi(a, \cdot) \in \mathfrak{S}_+$$

Faithful effect



=

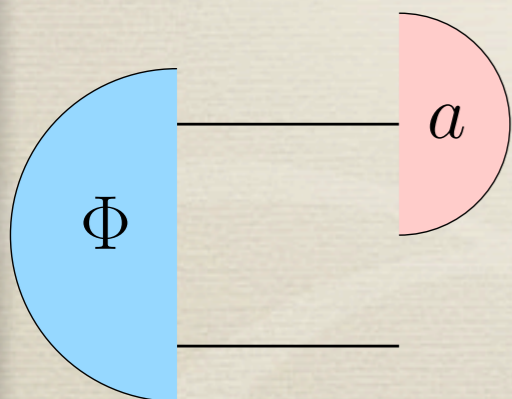


Remind the cone-isomorphism from the faithful state Φ

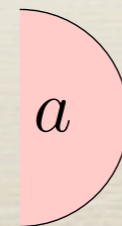
$$\mathfrak{E}_+ \ni a \mapsto \omega_a = \Phi(a, \cdot) \in \mathfrak{S}_+$$

FAITHFUL: There exist a bipartite effect F achieving the inverse of the isomorphism $a \mapsto \omega_a = \Phi(a, \cdot)$ namely:

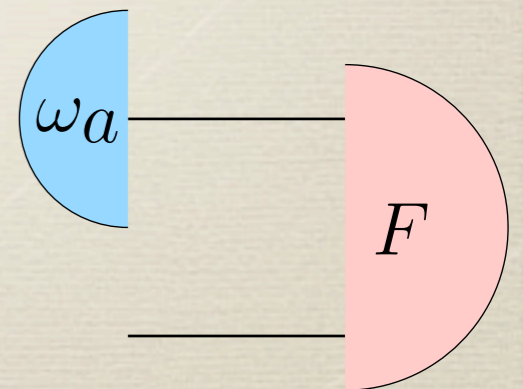
$$F_{23}(\omega_a)_2 = F_{23}\Phi_{12}(a, \cdot) = \alpha a_3, \quad 0 < \alpha \leq 1$$



=

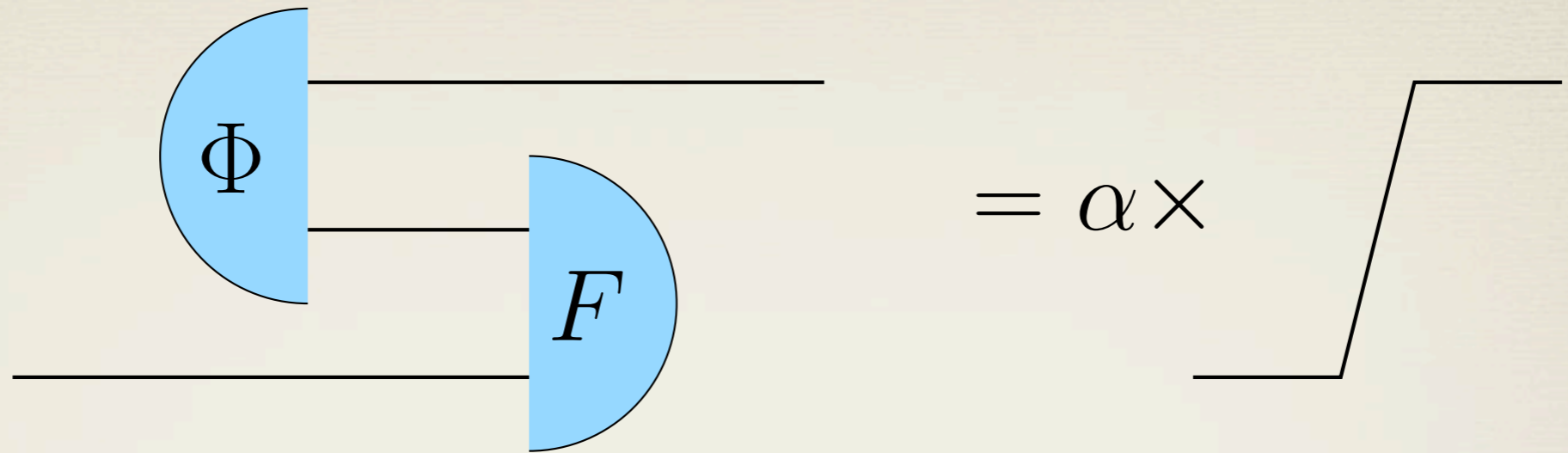


=



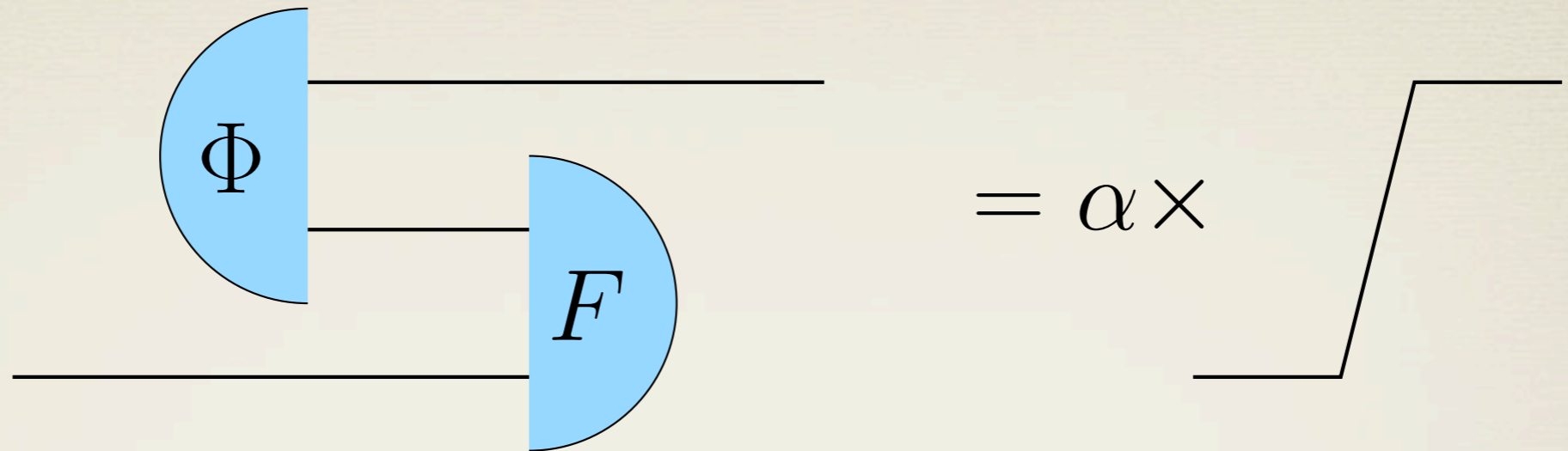
Consequences of FAITHE

► Teleportation:



Consequences of FAITHE

► Teleportation:



► F is **completely faithful**, i.e. $F_{\mathcal{A}} := F \circ (\mathcal{I}, \mathcal{A}) \iff \mathcal{A}$
 realizes the cone-isomorphism: $\mathfrak{E}_+(S^{\odot 2}) \simeq \mathfrak{I}_+(S)$

Consequences of FAITHE

▶ $\mathfrak{E}_+(\mathcal{S}^{\odot 2}) \ni A \mapsto \Omega_A := A_{23}(\Phi, \Phi) \in \mathfrak{S}_+(\mathcal{S}^{\odot 2})$

is a bijective map between $\mathfrak{S}_{\mathbb{F}}(\mathcal{S}^{\odot 2})$ and $\mathfrak{E}_{\mathbb{F}}(\mathcal{S}^{\odot 2})$

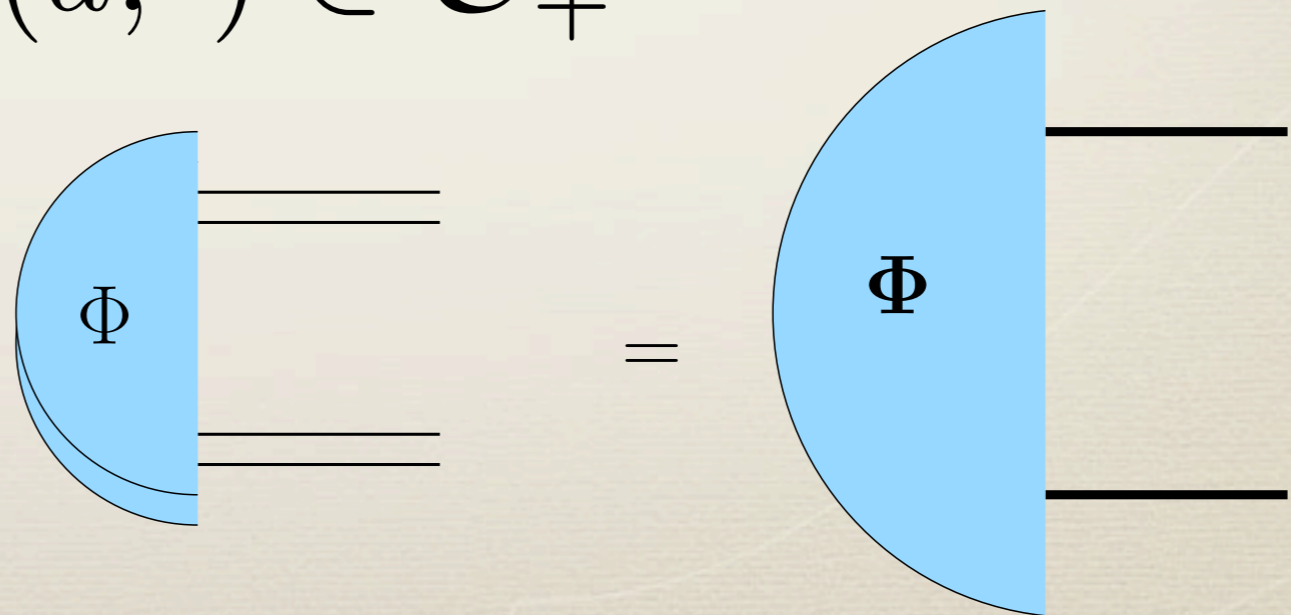
however, it does not necessarily realize the cone-isomorphism:

$$\mathfrak{S}_+(\mathcal{S}^{\odot 2}) \simeq \mathfrak{E}_+(\mathcal{S}^{\odot 2})$$

namely it is not the 4-partite equivalent of:

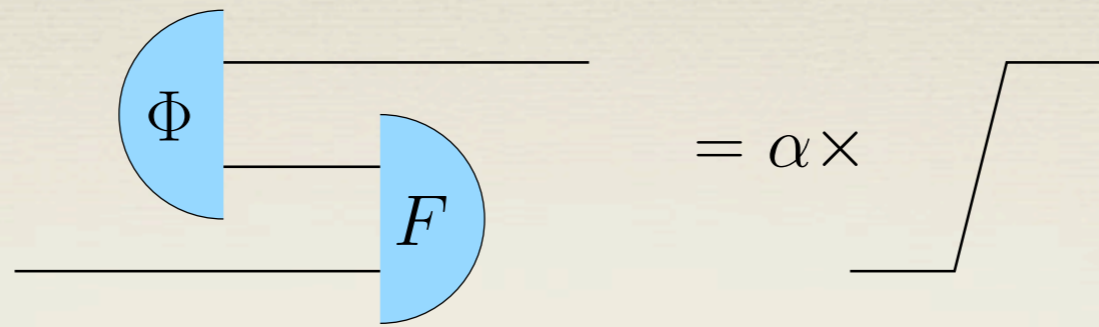
$$\mathfrak{E}_+ \ni a \mapsto \omega_a = \Phi(a, \cdot) \in \mathfrak{S}_+$$

i.e. we don't have:



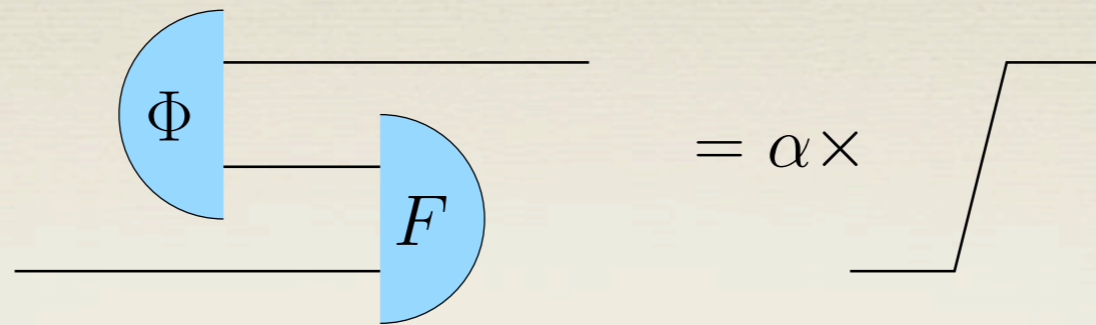
Consequences of FAITHE

Teleportation:



Consequences of FAITHE

Teleportation:

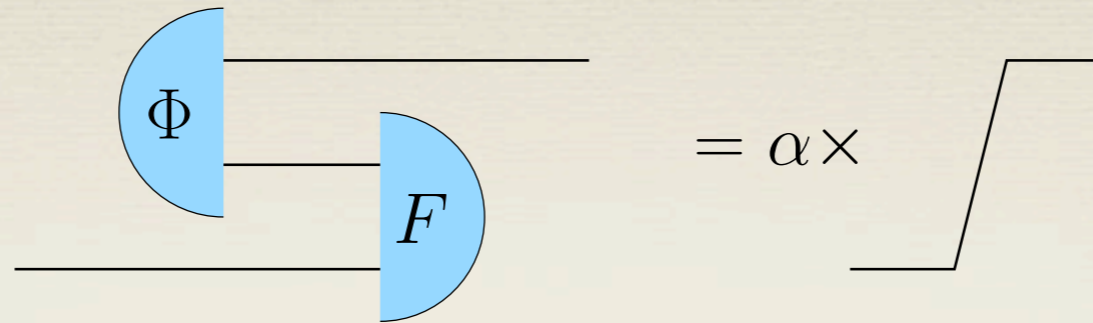


$$\alpha(\mathcal{S}) = \max_{E \in \mathfrak{E}(\mathcal{S}^{\odot 2})} \{(\Phi, \Phi)(e, E, e)\}$$

is a property of the system and depends on the particular probabilistic theory

Consequences of FAITHE

Teleportation:



$$\alpha(\mathcal{S}) = \max_{E \in \mathfrak{E}(\mathcal{S}^{\odot 2})} \{(\Phi, \Phi)(e, E, e)\}$$

is a property of the system and depends on the particular probabilistic theory

In Quantum Mechanics: $\alpha = \dim(\mathbf{H})^{-2}$

$$\omega_a = \sqrt{\alpha} \zeta(a)$$

$$(\cdot, F)(\Phi, \cdot) = \sqrt{\alpha} |\Phi|$$

Consequences of PURIFY

PURIFY: Every state has a purification on two identical systems.

Consequences of PURIFY

PURIFY: Every state has a purification on two identical systems.



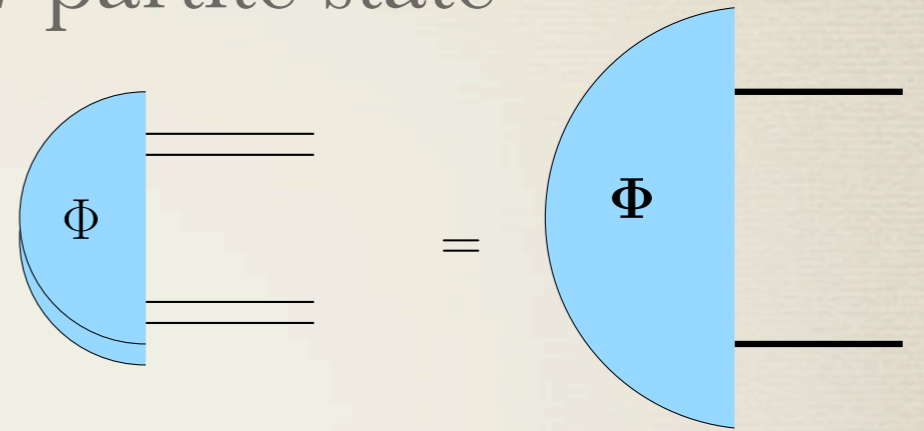
- ▶ Each state can be obtained by applying an atomic transformation to the marginal state $\chi = \Phi(e, \cdot)$
- ▶ Each effect contains an atomic transformation.
- ▶ \mathcal{I} is atomic.
- ▶ Φ is pure.

SUPER-PFAITH

SUPERFAITH: There exists a symmetric bipartite state such that for any number N of systems the $2N$ -partite state

$$\Phi_{134\dots 246\dots} := \Phi_{12}\Phi_{34}\Phi_{56} \dots$$

is preparationally faithful for $S^{\odot N}$

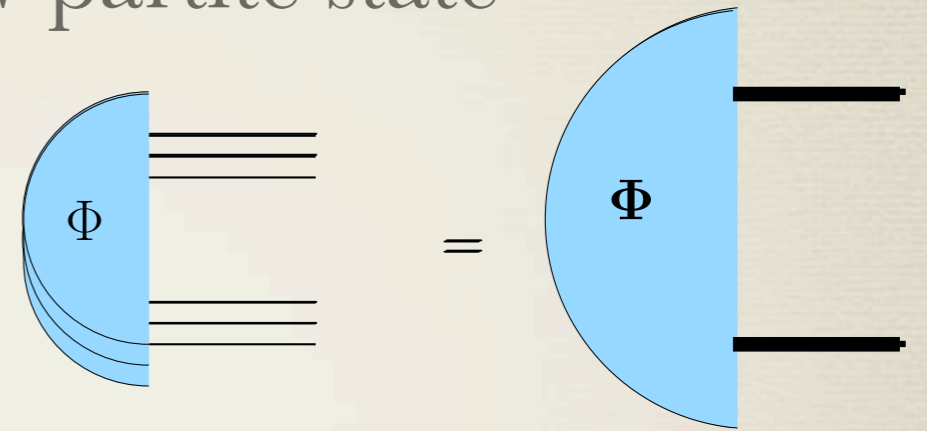


SUPER-PFAITH

SUPERFAITH: There exists a symmetric bipartite state such that for any number N of systems the $2N$ -partite state

$$\Phi_{134\dots 246\dots} := \Phi_{12}\Phi_{34}\Phi_{56} \dots$$

is preparationally faithful for $S^{\odot N}$

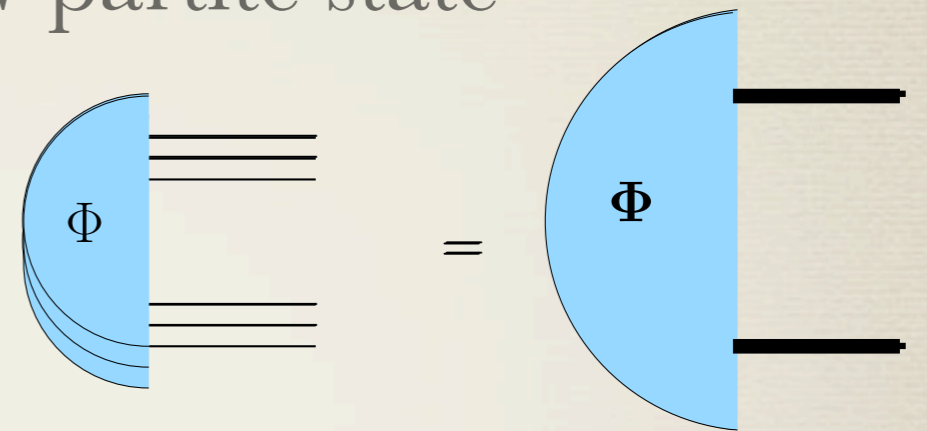


SUPER-PFAITH

SUPERFAITH: There exists a symmetric bipartite state such that for any number N of systems the $2N$ -partite state

$$\Phi_{134\dots 246\dots} := \Phi_{12}\Phi_{34}\Phi_{56} \dots$$

is preparationally faithful for $S^{\odot N}$



+PURIFY (with purifications connected by local automorphisms) \rightarrow
FAITHE + Stinespring:

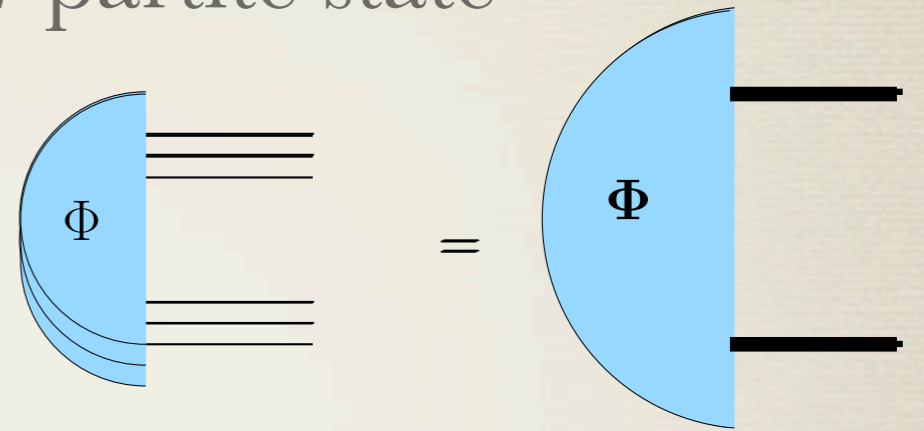
with G. Chiribella and
P. Perinotti

SUPER-PFAITH

SUPERFAITH: There exists a symmetric bipartite state such that for any number N of systems the $2N$ -partite state

$$\Phi_{134\dots 246\dots} := \Phi_{12}\Phi_{34}\Phi_{56} \dots$$

is preparationally faithful for $\mathcal{S}^{\odot N}$

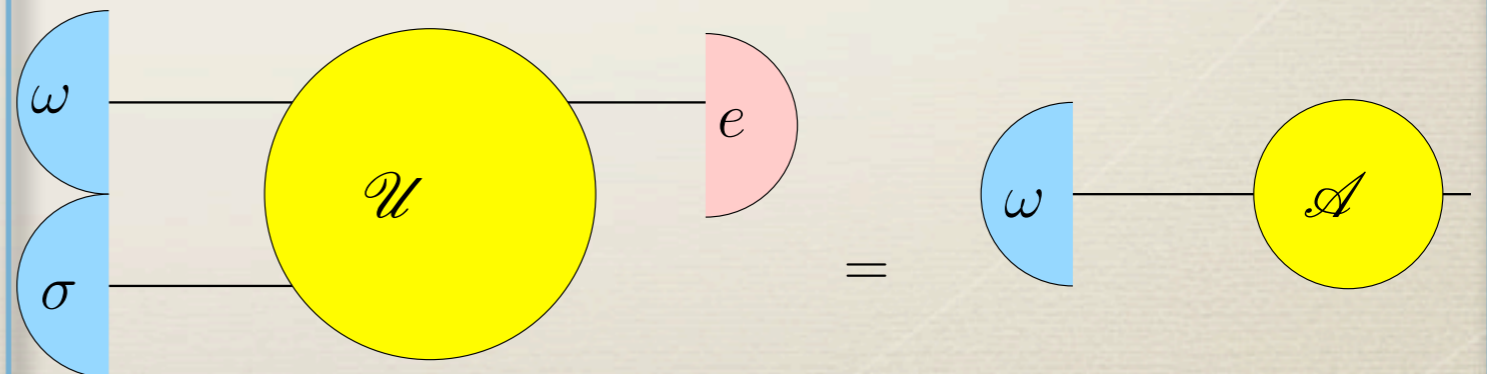


+PURIFY (with purifications connected by local automorphisms) \rightarrow

FAITHE+ Stinespring:

with G. Chiribella and P. Perinotti

$$\mathcal{U}_{12}(\sigma, \omega)(e, \cdot) = \mathcal{A}\omega$$



A teal background with a white speech bubble containing text. The speech bubble has a black outline and a tail pointing towards the bottom left. The text inside is in a teal serif font, with the first line in all caps and the second line in italics.

WHAT IS THE
SOMETHING MORE?

WHAT IS THE
SOMETHING MORE?

PFAITH
+FAITHE+PURIFY

WHAT IS THE
SOMETHING MORE?

OR

WHAT IS THE
SOMETHING MORE?

SUPER-PEFAITH
+PURIFY?

WHAT IS THE
SOMETHING MORE?

In any case, it must give that:
effects make a C^ -algebra*

RECONSTRUCTING QM FROM PROBABILITIES

Algebra of effects \Leftrightarrow Choi-Jamiołkowski isomorphism
+ atomicity of evolution

Quantum Tomography for Measuring Experimentally the Matrix Elements of an Arbitrary Quantum Operation

G. M. D'Ariano and P. Lo Presti

at our disposal a general method for experimentally determining the quantum operation matrix, using any available quantum-tomographic scheme for the system in consideration, and a single fixed state at the input, which is an entangled (not even maximally) state. In the optical domain we show that one can achieve the tomographic reconstruction of the operation using exactly the same apparatus of the recently performed experiment of Ref. [9].

Let us consider for simplicity a “pure” quantum operation in the form (5). Given an orthonormal basis $\{|j\rangle\}$ corresponding to some physical observable, how can we determine the matrix $A_{ij} = \langle i|A|j\rangle$ experimentally? Instead of acting with the contraction A on an “isolated” system, we perform the map on a system which is entangled in the state $|\psi\rangle\rangle \in \mathcal{H} \otimes \mathcal{H}$ with an identical system; i.e.,

$$|\psi\rangle\rangle \rightarrow |\phi\rangle\rangle = \frac{A \otimes I |\psi\rangle\rangle}{\|A\psi\|_{HS}}. \quad (6)$$

With the double ket we denote bipartite vectors $|\psi\rangle\rangle \in \mathcal{H} \otimes \mathcal{H}$, which, keeping the basis $\{|j\rangle\}$ as fixed, are in one-to-one correspondence with matrices as follows:

$$|\psi\rangle\rangle = \sum_{ij} \psi_{ij} |i\rangle \otimes |j\rangle. \quad (7)$$

$$A_{ij} = \kappa \langle E_{ij}(\psi) \rangle, \quad (10)$$

where the operator $E_{ij}(\psi)$ is given by

$$E_{ij}(\psi) = |i_0\rangle\langle i| \otimes |j_0\rangle\langle \psi^{-1*}(j)|, \quad (11)$$

and the proportionality constant is given by

$$\kappa = e^{i\theta} \sqrt{\frac{p_A(\psi)}{\langle |i_0, j_0\rangle\rangle \langle\langle i_0, j_0| \rangle\rangle}}. \quad (12)$$

Since A_{ij} is written only in terms of output ensemble averages, it can be estimated through quantum tomography. Quantum tomography [10,11] is a method to estimate the ensemble average $\langle H \rangle$ of any arbitrary operator H on \mathcal{H} by using only measurement outcomes of a *quorum* of observables $\{O(l)\}$. A *quorum* is just a set of operators $\{O(l)\}$ which are observable (i.e., have orthonormal resolution) and span the linear space of operators on \mathcal{H} . This means that any operator H can be expanded as $H = \sum_l \text{Tr}[Q^\dagger(l)H]O(l)$, where $\{Q(l)\}$ and $\{O(l)\}$ form a biorthogonal set such that $\text{Tr}[Q^\dagger(i)O(j)] = \delta_{ij}$. Hence, the tomographic estimation of the ensemble average $\langle H \rangle$ is obtained as the double average—over both the ensemble and the quorum—of the unbiased

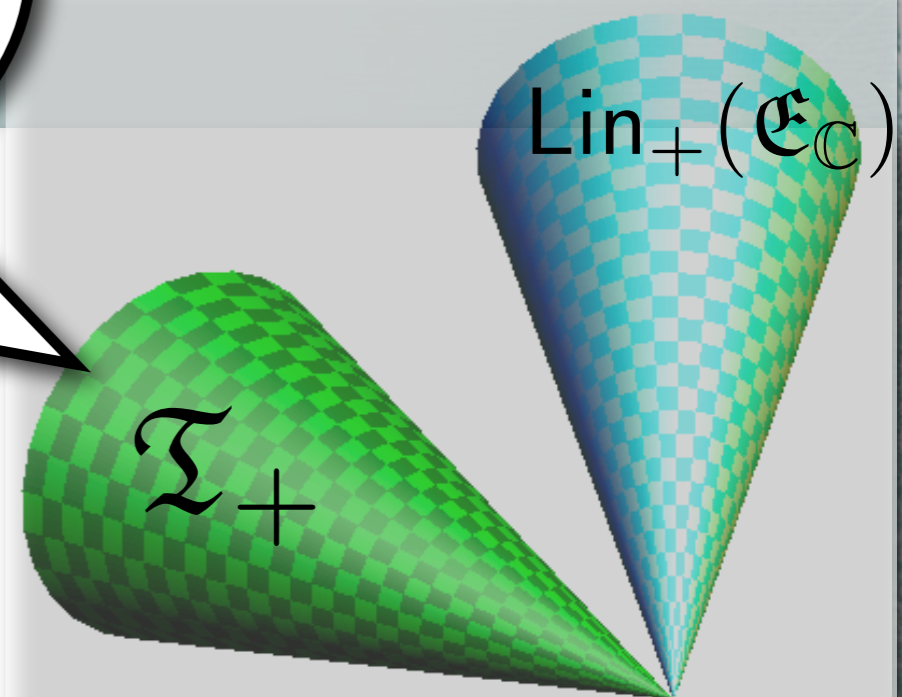
A teal background with a white speech bubble containing text. The speech bubble has a black outline and a tail pointing towards the bottom left. The text inside is in a dark teal, serif font, centered within the bubble.

DIFFICULT PROBLEM:
HOW TO INTRODUCE
COMPOSITION OF
EFFECTS?

**DIFFICULT PROBLEM:
HOW TO INTRODUCE
COMPOSITION OF
EFFECTS?**

**Equivalent to:
Choi-Jamiołkowski
isomorphism**

$$\mathbf{I} : \mathfrak{L}_+ \simeq \text{Lin}_+(\mathcal{E}_{\mathbb{C}})$$

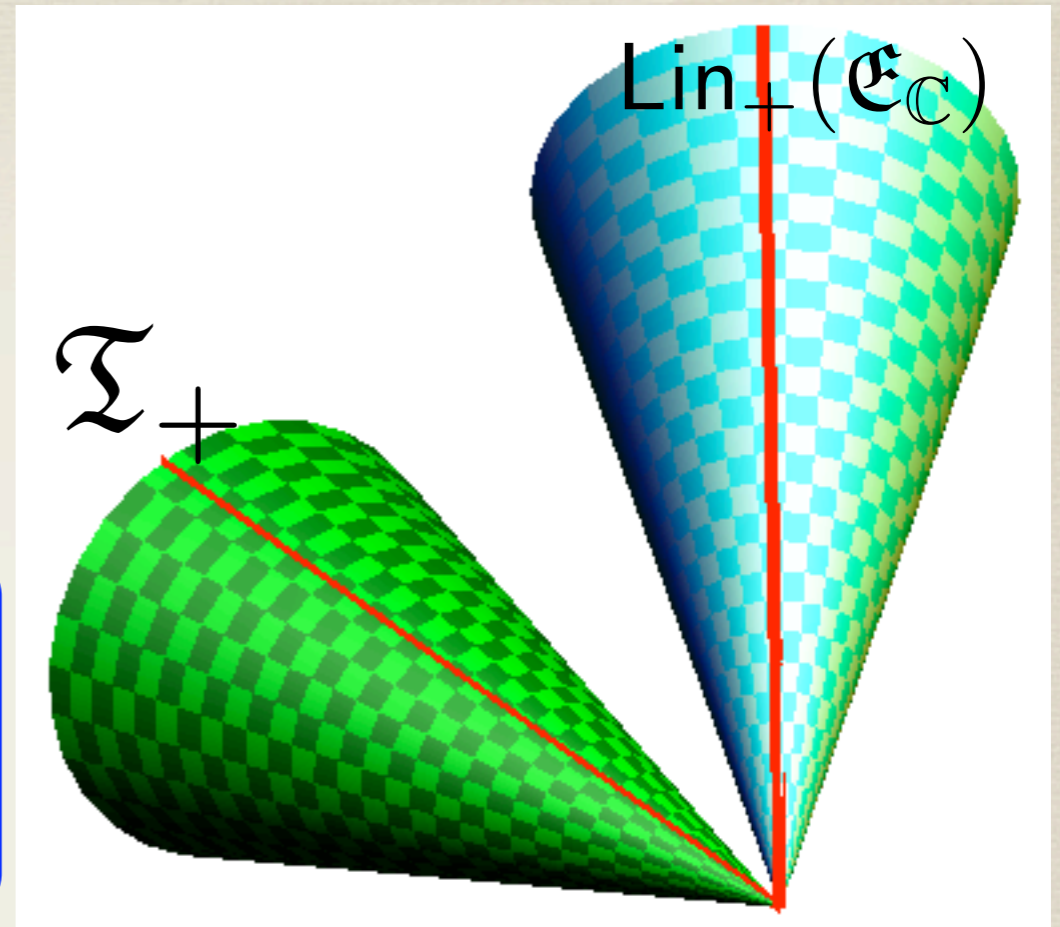


CJ Isomorphism \Rightarrow composition of effects

Effects are identified with “atomic” events

(apart from a phase) i.e. events that cannot be written as sum of other events

AE (Atomicity of evolution):
the composition of “atomic” events is atomic

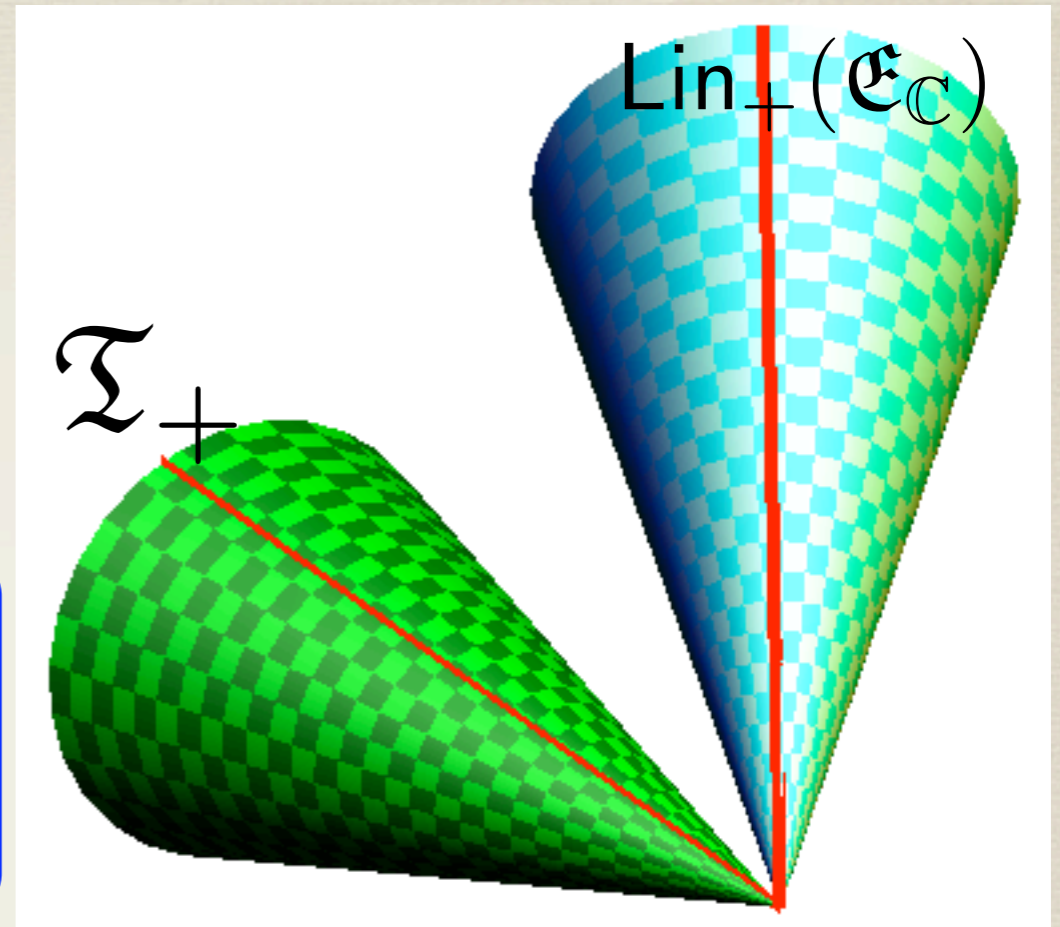


CJ Isomorphism \Rightarrow composition of effects

Effects are identified with “atomic” events

(apart from a phase) i.e. events that cannot be written as sum of other events

AE (Atomicity of evolution):
the composition of “atomic” events is atomic



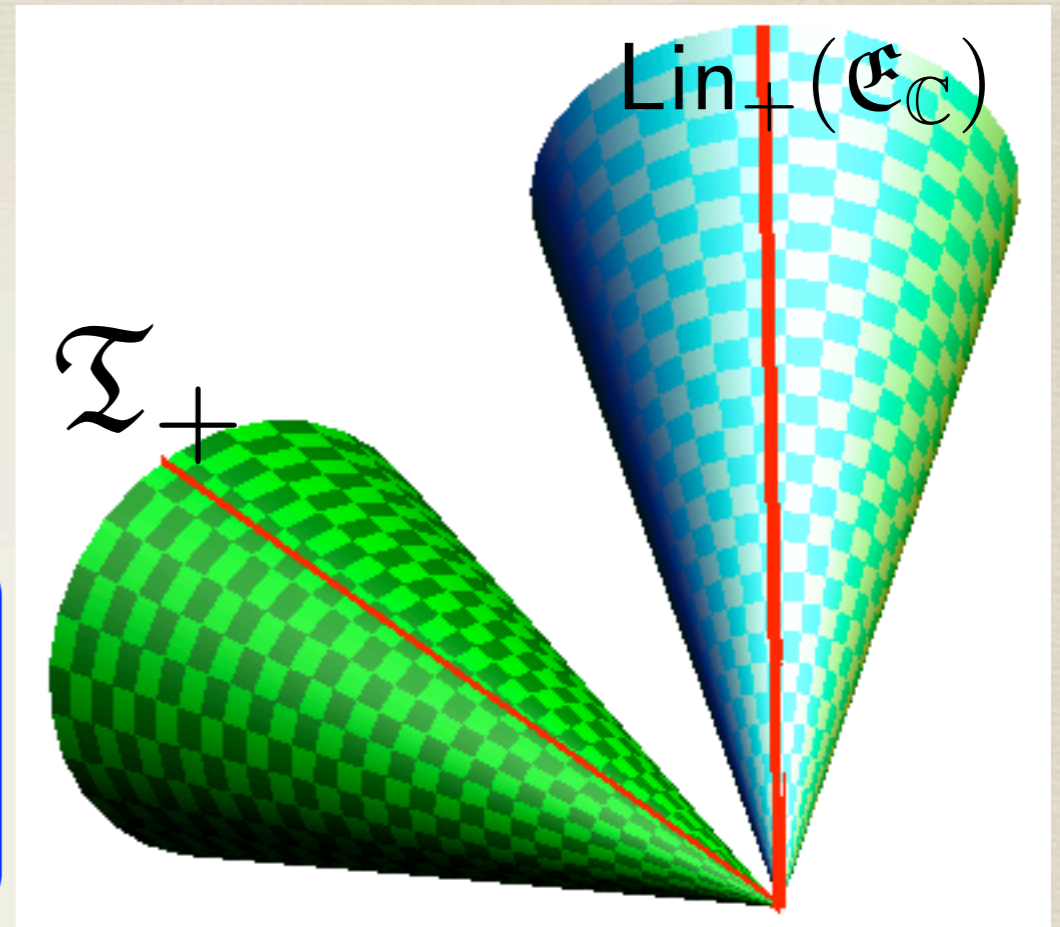
One can prove that the phase (two-cocycle) is trivial.

CJ Isomorphism \Rightarrow composition of effects

Effects are identified with “atomic” events

(apart from a phase) i.e. events that cannot be written as sum of other events

AE (Atomicity of evolution):
the composition of “atomic” events is atomic



One can prove that the phase (two-cocycle) is trivial.

Introduce the generalized transformation via the polarization identity:

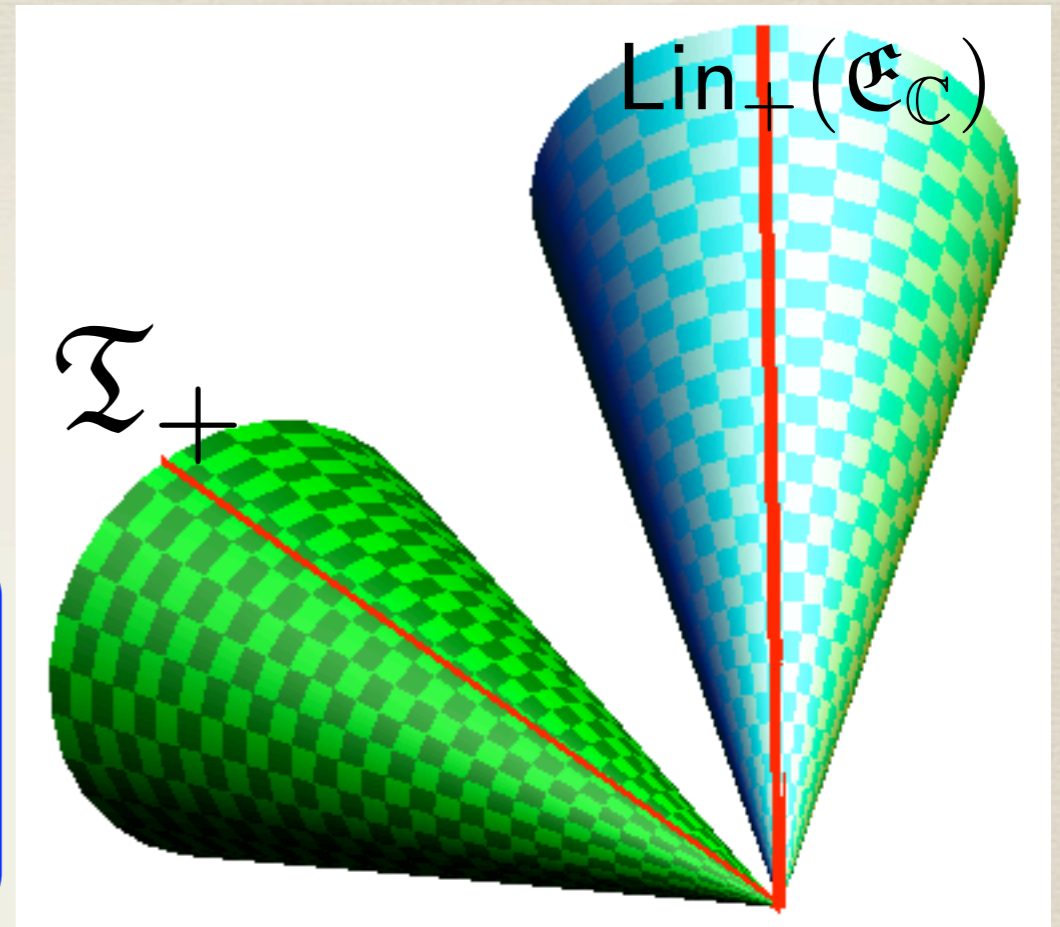
$$\mathcal{T}_{a,b} := \frac{1}{4} \sum_{k=0}^3 i^k \mathcal{T}_{a+i^k b}$$

CJ Isomorphism \Rightarrow composition of effects

Effects are identified with “atomic” events

(apart from a phase) i.e. events that cannot be written as sum of other events

AE (Atomicity of evolution): the composition of “atomic” events is atomic



One can prove that the phase (two-cocycle) is trivial.
Introduce the generalized transformation via the polarization identity:

$$\mathcal{T}_{a,b} := \frac{1}{4} \sum_{k=0}^3 i^k \mathcal{T}_{a+i^k b}$$

composition of effects as: $ab = e \circ \mathcal{T}_{e,a} \circ \mathcal{T}_{e,b}$

SUMMARY

SUMMARY

NSF

events \equiv transformations

evolution \equiv conditioning

SUMMARY

NSF

events \equiv transformations

evolution \equiv conditioning

Probabilistic
framework

C*-algebra of
transformations

independent systems
 \equiv no-signaling

SUMMARY

NSF

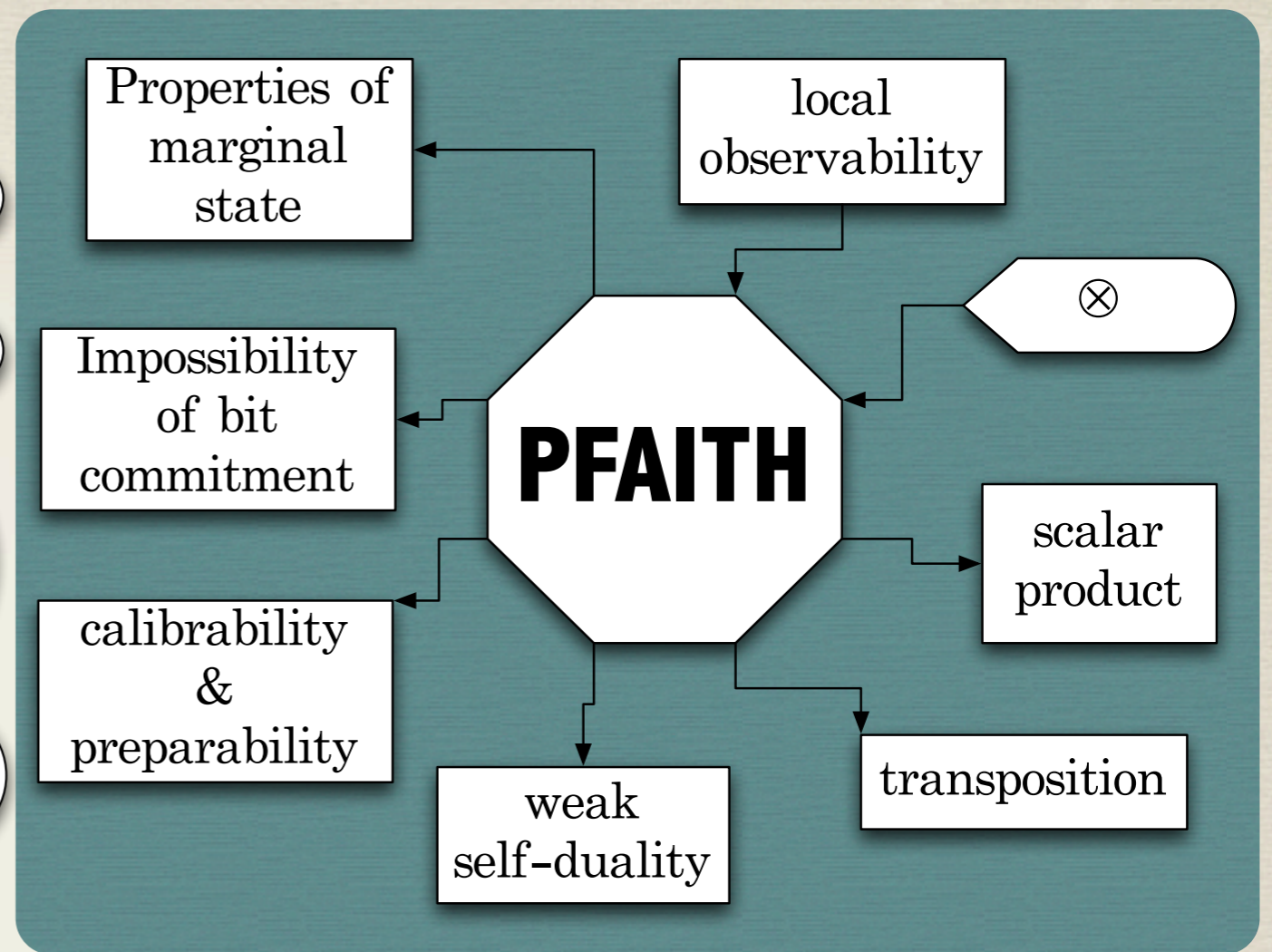
events \equiv transformations

evolution \equiv conditioning

Probabilistic framework

C*-algebra of transformations

independent systems \equiv no-signaling



SUMMARY

NSF

events \equiv transformations

evolution \equiv conditioning

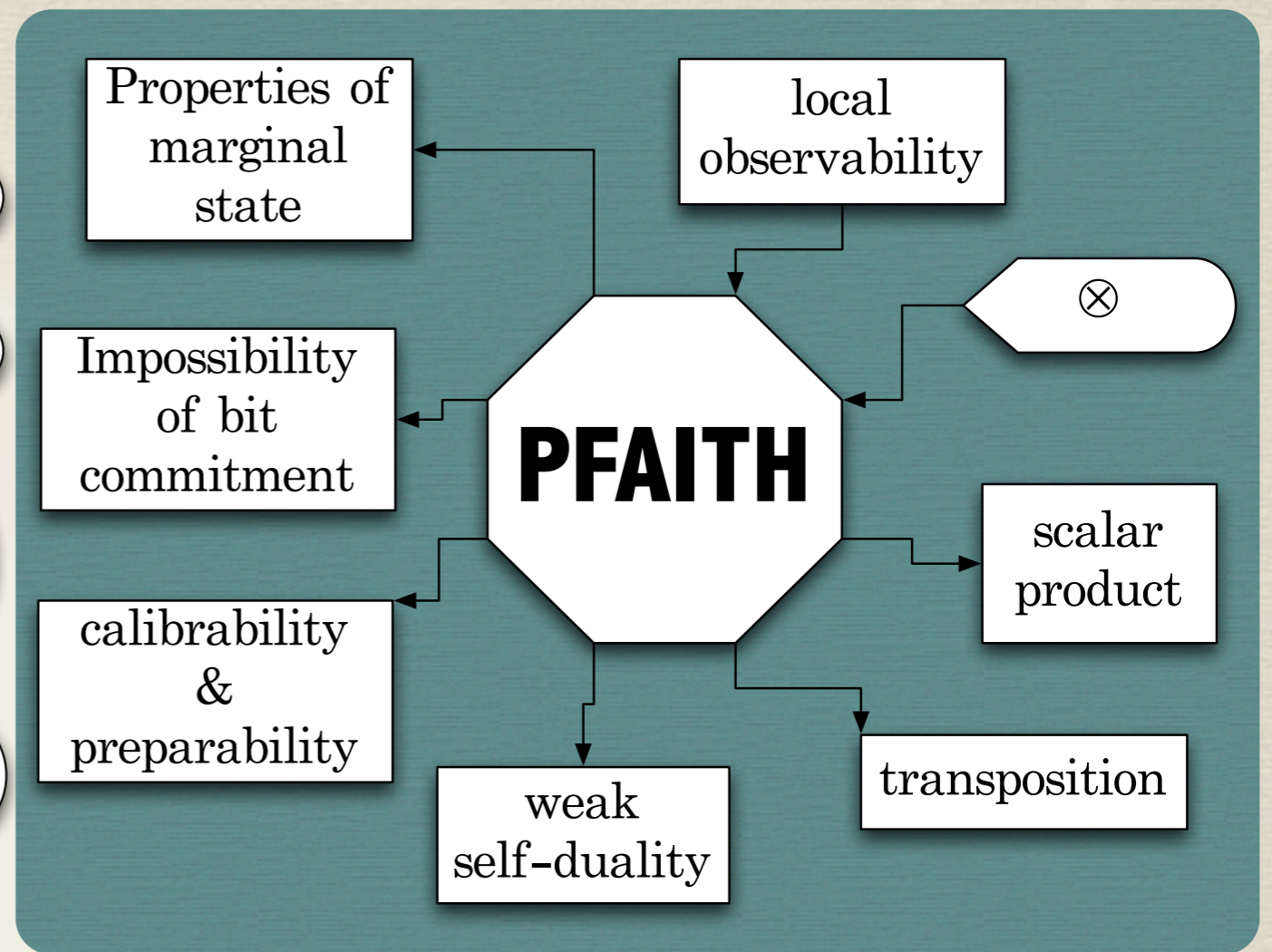
Probabilistic framework

C*-algebra of transformations

independent systems \equiv no-signaling

+FAITHE

TELEPORTATION



+PURIFY

States are orbit of atomic transformations on the marginal state

\mathcal{I} is atomic

Φ is pure

Each effect contains an atomic transformation

SUMMARY

